FAUST Quick Reference

(version 2.15.0)

GRAME Centre National de Création Musicale

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Chapter 1

Introduction

FAUST (Functional Audio Stream) is a functional programming language specifically designed for real-time signal processing and synthesis. FAUST targets high-performance signal processing applications and audio plug-ins for a variety of platforms and standards.

1.1 Design Principles

Various principles have guided the design of FAUST:

- FAUST is a *specification language*. It aims at providing an adequate notation to describe *signal processors* from a mathematical point of view. FAUST is, as much as possible, free from implementation details.
- FAUST programs are fully compiled, not interpreted. The compiler translates
 FAUST programs into equivalent C++ programs taking care of generating the
 most efficient code. The result can generally compete with, and sometimes even
 outperform, C++ code written by seasoned programmers.
- The generated code works at the sample level. It is therefore suited to implement low-level DSP functions like recursive filters. Moreover the code can be easily embedded. It is self-contained and doesn't depend of any DSP library or runtime system. It has a very deterministic behavior and a constant memory footprint.
- The semantic of FAUST is simple and well defined. This is not just of academic
 interest. It allows the FAUST compiler to be semantically driven. Instead of
 compiling a program literally, it compiles the mathematical function it denotes.
 This feature is useful for example to promote components reuse while preserving
 optimal performance.

- FAUST is a textual language but nevertheless block-diagram oriented. It actually combines two approaches: *functional programming* and *algebraic block-diagrams*. The key idea is to view block-diagram construction as function composition. For that purpose, FAUST relies on a *block-diagram algebra* of five composition operations (: , ~ <: :>).
- Thanks to the notion of *architecture*, FAUST programs can be easily deployed on a large variety of audio platforms and plugin formats without any change to the FAUST code.

1.2 Signal Processor Semantic

A FAUST program describes a *signal processor*. The role of a *signal processor* is to transforms a (possibly empty) group of *input signals* in order to produce a (possibly empty) group of *output signals*. Most audio equipments can be modeled as *signal processors*. They have audio inputs, audio outputs as well as control signals interfaced with sliders, knobs, vu-meters, etc.

More precisely:

- viole precisely.
 - A signal s is a discrete function of time s: Z → R. The value of a signal s at time t is written s(t). The values of signals are usually needed starting from time 0. But to take into account delay operations, negative times are possible and are always mapped to zeros. Therefore for any FAUST signal s we have ∀t < 0, s(t) = 0. In operational terms this corresponds to assuming that all delay lines are signals initialized with 0s.
 - The set of all possible signals is $\mathbb{S} = \mathbb{Z} \to \mathbb{R}$.
 - A group of n signals (a n-tuple of signals) is written $(s_1, \ldots, s_n) \in \mathbb{S}^n$. The *empty tuple*, single element of \mathbb{S}^0 is notated ().
 - A signal processors p, is a function from n-tuples of signals to m-tuples of signals $p: \mathbb{S}^n \to \mathbb{S}^m$. The set $\mathbb{P} = \bigcup_{n,m} \mathbb{S}^n \to \mathbb{S}^m$ is the set of all possible signal processors.

As an example, let's express the semantic of the Faust primitive +. Like any Faust expression, it is a signal processor. Its signature is $\mathbb{S}^2 \to \mathbb{S}$. It takes two input signals X_0 and X_1 and produce an output signal Y such that $Y(t) = X_0(t) + X_1(t)$.

Numbers are signal processors too. For example the number 3 has signature $\mathbb{S}^0 \to \mathbb{S}$. It takes no input signals and produce an output signal Y such that Y(t) = 3.

Faust considers two type of signals: integer signals $(s: \mathbb{Z} \to \mathbb{Z})$ and floating point signals $(s: \mathbb{Z} \to \mathbb{Q})$ Exchanges with the outside world are, by convention, made using floating point signals. The full range is represented by sample values between -1.0 and +1.0.

Chapter 2

Compiling and installing FAUST

The FAUST source distribution faust-2.15.0.tar.gz can be downloaded from GitHub (https://github.com/grame-cncm/faust/releases).

2.1 Organization of the distribution

The first thing is to decompress the downloaded archive.

```
tar xzf faust-2.15.0.tar.gz
```

The resulting faust-2.15.0/ folder should contain the following elements:

architecture/ FAUST libraries and architecture files

benchmark tools to measure the efficiency of the generated code

compiler/ sources of the FAUST compiler examples/ examples of FAUST programs

syntax-highlighting/ support for syntax highlighting for several editors documentation/ FAUST's documentation, including this manual tools/ tools to produce audio applications and plugins

COPYING license information

Makefile used to build and install FAUST
README instructions on how to build and install FAUST

2.2 Compilation

Faust has no dependencies outside standard libraries. Therefore the compilation should be straightforward. There is no configuration phase, to compile the Faust compiler simply do:

```
cd faust-2.15.0/
make
```

If the compilation was successful you can test the compiler before installing it:

```
[cd faust-2.15.0/]
./compiler/faust -v
```

It should output:

```
FAUST, DSP to C++ compiler, Version 2.15.0
Copyright (C) 2002-2019, GRAME - Centre...
```

Then you can also try to compile one of the examples:

```
[cd faust-2.15.0/]
./compiler/faust examples/generator/noise.dsp
```

It should produce some C++ code on the standard output

2.3 Installation

You can install FAUST with:

```
[cd faust-2.15.0/]
sudo make install
```

or

```
[cd faust-2.15.0/]
su
make install
```

depending on your system.

2.4 Compilation of the examples

Once Faust correctly installed, you can have a look at the provided examples in the examples/ folder. This folder contains a Makefile with all the required instructions to build these examples for various *architectures*, either standalone audio applications or plugins.

The command make help will list the available targets. Before using a specific target, make sure you have the appropriate development tools, libraries and headers installed. For example to compile the examples as ALSA applications with a GTK user interface do a make alsagtk. This will create a alsagtkdir/subfolder with all the binaries.

An architecture file provides the code needed to connect a signal processor to the outside world. It typically defines the audio communications and user interface.

Chapter 3

Faust syntax

This section describes the syntax of FAUST. Figure 3.1 gives an overview of the various concepts and where they are defined in this section.

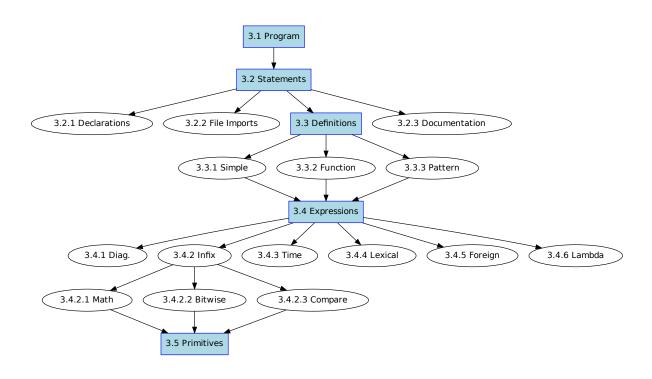


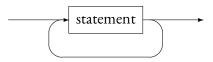
Figure 3.1: Overview of FAUST syntax

As we will see, definitions and expressions have a central role.

3.1 FAUST program

A FAUST program is essentially a list of *statements*. These statements can be *declarations*, *imports*, *definitions* and *documentation tags*, with optional C++ style (//... and /*...*/) comments.

program



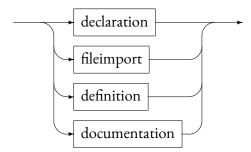
Here is a short FAUST program that implements of a simple noise generator. It exhibits various kind of statements: two *declarations*, an *import*, a *comment* and a *definition*. We will see later on *documentation* statements (3.2.3).

The keyword process is the equivalent of main in C/C++. Any FAUST program, to be valid, must at least define process.

3.2 Statements

The *statements* of a FAUST program are of four kinds: *metadata declarations*, *file imports*, *definitions* and *documentation*. All statements but documentation end with a semicolon (;).

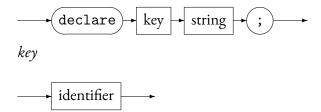
statement



3.2.1 Declarations

Meta-data declarations (for example declare name "noise";) are optional and typically used to document a FAUST project.

declaration



Contrary to regular comments, these declarations will appear in the C++ code generated by the compiler. A good practice is to start a FAUST program with some standard declarations:

```
declare name "MyProgram";
declare author "MySelf";
declare copyright "MyCompany";
declare version "1.00";
declare license "BSD";
```

3.2.2 Imports

File imports allow to import definitions from other source files.

fileimport



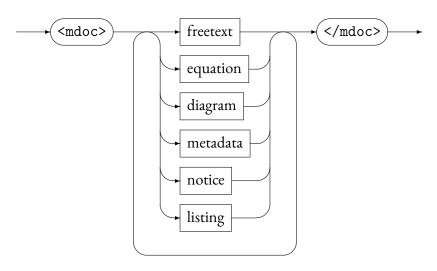
For example import("math.lib"); imports the definitions of the math.lib library, a set of additional mathematical functions provided as foreign functions.

3.2.3 Documentation

Documentation statements are optional and typically used to control the generation of the mathematical documentation of a FAUST program. This documentation system is detailed chapter 12. In this section we will essentially describe the documentation statements syntax.

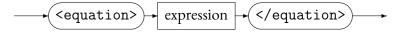
A documentation statement starts with an opening <mdoc> tag and ends with a closing </mdoc> tag. Free text content, typically in LaTeX format, can be placed in between these two tags.

documentation



Moreover, optional sub-tags can be inserted in the text content itself to require the generation, at the insertion point, of mathematical *equations*, graphical *block-diagrams*, FAUST source code *listing* and explanation *notice*.

equation



The generation of the mathematical equations of a Faust expression can be requested by placing this expression between an opening <equation> and a closing </equation> tag. The expression is evaluated within the lexical context of the Faust program.

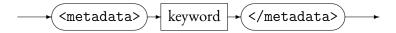
diagram



Similarly, the generation of the graphical block-diagram of a Faust expression can be requested by placing this expression between an opening <diagram> and a closing </diagram> tag. The expression is evaluated within the lexical context of the Faust program.

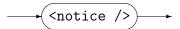
17

metadata



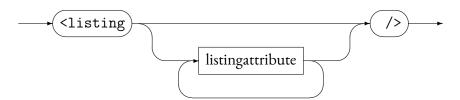
The <metadata> tags allow to reference FAUST metadatas (cf. declarations), calling the corresponding keyword.

notice

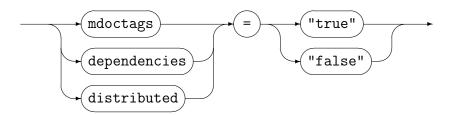


The <notice /> empty-element tag is used to generate the conventions used in the mathematical equations.

listing



listingattribute



The sting /> empty-element tag is used to generate the listing of the FAUST program. Its three attributes mdoctags, dependencies and distributed enable or disable respectively <mdoc> tags, other files dependencies and distribution of interleaved faust code between <mdoc> sections.

3.3 Definitions

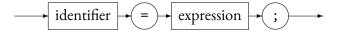
A *definition* associates an identifier with an expression it stands for.

Definitions are essentially a convenient shortcut avoiding to type long expressions. During compilation, more precisely during the evaluation stage, identifiers are replaced by their definitions. It is therefore always equivalent to use an identifier or directly its definition. Please note that multiple definitions of a same identifier are not allowed, unless it is a pattern matching based definition.

3.3.1 Simple Definitions

The syntax of a simple definition is:

definition



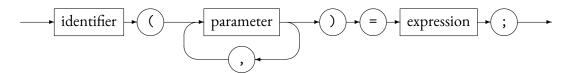
For example here is the definition of random, a simple pseudo-random number generator:

```
random = +(12345) ~ *(1103515245);
```

3.3.2 Function Definitions

Definitions with formal parameters correspond to functions definitions.

definition



For example the definition of linear2db, a function that converts linear values to decibels, is:

```
linear2db(x) = 20*log10(x);
```

Please note that this notation is only a convenient alternative to the direct use of *lambda-abstractions* (also called anonymous functions). The following is an equivalent definition of linear2db using a lambda-abstraction:

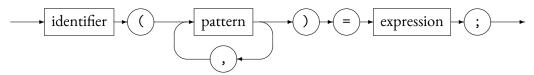
```
linear2db = (x).(20*log10(x));
```

19

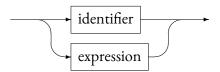
3.3.3 Definitions with pattern matching

Moreover, formal parameters can also be full expressions representing patterns.

definition



pattern



This powerful mechanism allows to algorithmically create and manipulate block diagrams expressions. Let's say that you want to describe a function to duplicate an expression several times in parallel:

```
duplicate(1,x) = x;
duplicate(n,x) = x, duplicate(n-1,x);
```

Please note that this last definition is a convenient alternative to the more verbose:

Here is another example to count the number of elements of a list. Please note that we simulate lists using parallel composition: (1,2,3,5,7,11). The main limitation of this approach is that there is no empty list. Moreover lists of only one element are represented by this element:

```
count((x,xs)) = 1+count(xs);
count(x) = 1;
```

If we now write count (duplicate (10,666)) the expression will be evaluated to 10.

Please note that the order of pattern matching rules matters. The more specific rules must precede the more general rules. When this order is not respected, as in :

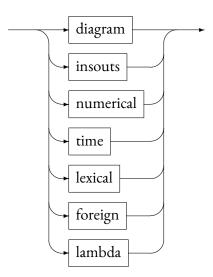
```
count(x) = 1;
count((x,xs)) = 1+count(xs);
```

the first rule will always match and the second rule will never be called.

3.4 Expressions

Despite its textual syntax, Faust is conceptually a block-diagram language. Faust expressions represent DSP block-diagrams and are assembled from primitive ones using various *composition* operations. More traditional *numerical* expressions in infix notation are also possible. Additionally Faust provides time based expressions, like delays, expressions related to lexical environments, expressions to interface with foreign function and lambda expressions.

expression



3.4.1 Diagram Expressions

Diagram expressions are assembled from primitive ones using either binary composition operations or high level iterative constructions.

diagramexp

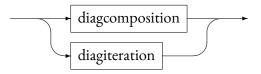
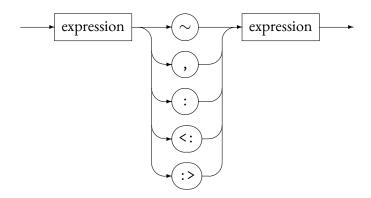


Diagram composition operations

Five binary *composition operations* are available to combine block-diagrams: *recursion*, *parallel*, *sequential*, *split* and *merge* composition. One can think of each of these composition operations as a particular way to connect two block diagrams.

diagcomposition



To describe precisely how these connections are done, we have to introduce some notation. The number of inputs and outputs of a bloc-diagram A are notated inputs (A) and outputs (A). The inputs and outputs themselves are respectively notated : [0]A, [1]A, [2]A, . . . and A[0], A[1], A[2], etc..

For each composition operation between two block-diagrams A and B we will describe the connections $A[i] \rightarrow [j]B$ that are created and the constraints on their relative numbers of inputs and outputs.

The priority and associativity of this five operations are given table 3.1.

Syntax	Pri.	Assoc.	Description
expression ~ expression	4	left	recursive composition
expression, expression	3	right	parallel composition
expression: expression	2	right	sequential composition
expression <: expression	I	right	split composition
expression :> expression	I	right	merge composition

Table 3.1: Block-Diagram composition operation priorities

Parallel Composition The *parallel composition* (A,B) (figure 3.2) is probably the simplest one. It places the two block-diagrams one on top of the other, without connections. The inputs of the resulting block-diagram are the inputs of A and B. The outputs of the resulting block-diagram are the outputs of A and B.

Parallel composition is an associative operation: (A, (B,C)) and ((A,B),C) are equivalents. When no parenthesis are used: A,B,C,D, FAUST uses right associativity and therefore build internally the expression (A, (B, (C,D))). This organization is important to know when using pattern matching techniques on parallel compositions.

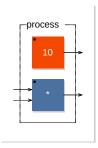


Figure 3.2: Example of parallel composition (10,*)

Sequential Composition The *sequential composition A*:B (figure 3.3) expects:

$$outputs(A) = inputs(B)$$
 (3.1)

It connects each output of *A* to the corresponding input of *B*:

$$A[i] \to [i]B \tag{3.2}$$

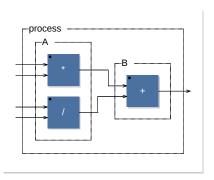


Figure 3.3: Example of sequential composition ((*,/):+)

Sequential composition is an associative operation: (A:(B:C)) and ((A:B):C) are equivalents. When no parenthesis are used, like in A:B:C:D, FAUST uses right associativity and therefore build internally the expression (A:(B:(C:D))).

Split Composition The *split composition* $A \le B$ (figure 3.4) operator is used to distribute the outputs of A to the inputs of B.

For the operation to be valid the number of inputs of B must be a multiple of the number of outputs of A:

$$outputs(A).k = inputs(B)$$
(3.3)

Each input i of B is connected to the output $i \mod k$ of A:

$$A[i \bmod k] \to [i]B \tag{3.4}$$

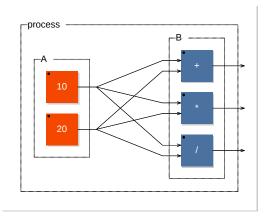


Figure 3.4: example of split composition ((10,20)<: (+,*,/))

Merge Composition The *merge composition* A:>B (figure 3.5) is the dual of the *split composition*. The number of outputs of A must be a multiple of the number of inputs of B:

$$outputs(A) = k.inputs(B)$$
 (3.5)

Each output i of A is connected to the input $i \mod k$ of B:

$$A[i] \to [i \bmod k]B \tag{3.6}$$

The k incoming signals of an input of B are summed together.

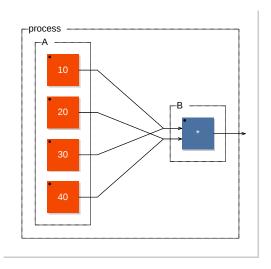


Figure 3.5: example of merge composition ((10,20,30,40):> *)

Recursive Composition The *recursive composition* A~B (figure 3.6) is used to create cycles in the block-diagram in order to express recursive computations. It is the most complex operation in terms of connections.

To be applicable it requires that:

$$\operatorname{outputs}(A) \ge \operatorname{inputs}(B) \operatorname{andinputs}(A) \ge \operatorname{outputs}(B)$$
 (3.7)

Each input of B is connected to the corresponding output of A via an implicit 1-sample delay :

$$A[i] \stackrel{Z^{-1}}{\to} [i]B \tag{3.8}$$

and each output of *B* is connected to the corresponding input of *A*:

$$B[i] \to [i]A \tag{3.9}$$

The inputs of the resulting block diagram are the remaining unconnected inputs of A. The outputs are all the outputs of A.

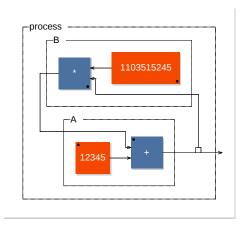
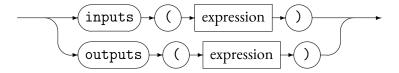


Figure 3.6: example of recursive composition + (12345) * *(1103515245)

Inputs and outputs of an expression

These two constructions can be used to know at compile time the number of inputs and outputs of any Faust expression.

insouts



They are useful to define high order functions and build algorithmically complex block-diagrams. Here is an example to automatically reverse the order of the outputs of an expression.

```
Xo(expr) = expr <: par(i,n,selector(n-i-1,n))
with { n=outputs(expr); };</pre>
```

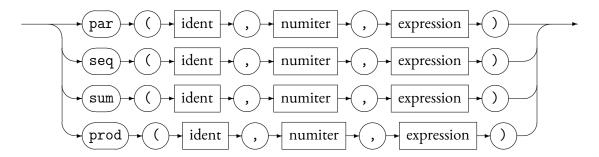
And the inputs of an expression:

For example Xi(-) will reverse the order of the two inputs of the substraction.

Iterations

Iterations are analogous to for (...) loops and provide a convenient way to automate some complex block-diagram constructions.

diagiteration



The following example shows the usage of seq to create a 10-bands filter:

numiter

```
expression
```

The number of iterations must be a constant expression.

3.4.2 Infix Notation and Other Syntax Extensions

Infix notation is commonly used in mathematics. It consists in placing the operand between the arguments as in 2+3

Besides its algebra-based core syntax, Faust provides some syntax extensions, in particular the familiar *infix notation*. For example if you want to multiply two numbers, let say 2 and 3, you can write directly 2*3 instead of the equivalent core-syntax expression 2,3:*.

The *infix notation* is not limited to numbers or numerical expressions. Arbitrary expressions A and B can be used, provided that A, B has exactly two outputs. For example you can write _/2. It is equivalent to _,2:/ which divides the incoming signal by 2.

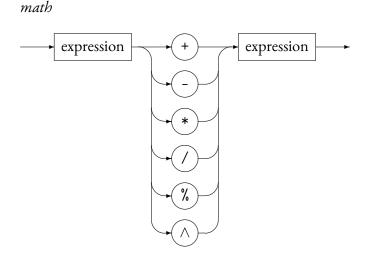
Examples of equivalences are given table 3.2.

Table 3.2: Infix and core syntax equivalences

In case of doubts on the meaning of an infix expression, for example _*_, it is useful to translate it to its core syntax equivalent, here _ , _ : *, which is equivalent to *.

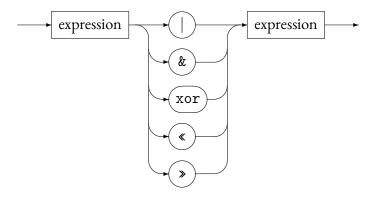
Infix Operators

Built-in primitives that can be used in infix notation are called *infix operators* and are listed in this section. Please note that a more detailed description of these operators is available section 3.5.

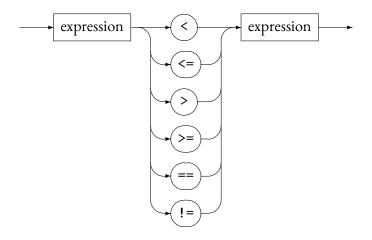


27

bitwise



comparison



delay



Prefix Notation

Beside *infix notation*, it is also possible to use *prefix notation*. The *prefix notation* is the usual mathematical notation for functions f(x, y, z, ...), but extended to *infix operators*.

It consists in first having the operator, for example /, followed by its arguments between parentheses: /(2,3) (see table 3.3).

```
*(2,3) \equiv 2,3 : *
@(\_,7) \equiv \_,7 : @
/(\_,2) \equiv \_,2 : /
<(A,B) \equiv A,B : <
```

Table 3.3: Prefix to core syntax translation rules

Partial Application

The partial application notation is a variant of prefix notation in which not all arguments are given. For instance /(2) (divide by 2), ^(3) (rise to the cube) and @(512) (delay by 512 samples) are examples of partial applications where only one argument is given. The result of a partial application is a function that "waits" for the remaining arguments.

When doing partial application with an *infix operator*, it is important to note that the supplied argument is not the first argument, but always the second one, as summarized table 3.4.

```
+(C) = _,C:*

-(C) = _,C:-

<(C) = _,C:<

/(C) = _,C:/
```

Table 3.4: Partial application of infix operators

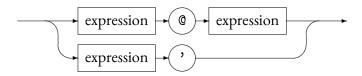
For commutative operations that doesn't matter. But for non-commutative ones, it is more "natural" to fix the second argument. We use divide by 2(/(2)) or rise to the cube $(^{(3)})$ more often than the other way around.

Please note that this rule only applies to infix operators, not other primitives or functions. If you partially apply a regular function to a single argument, it will correspond to the first parameter.

3.4.3 Time expressions

Time expressions are used to express delays. The notation X@10 represent the signal X delayed by 10 samples. The notation X' represent the signal X delayed by one sample and is therefore equivalent to X@1.

time



The delay don't have to be fixed, but it must be positive and bounded. The values of a slider are perfectly acceptable as in the following example:

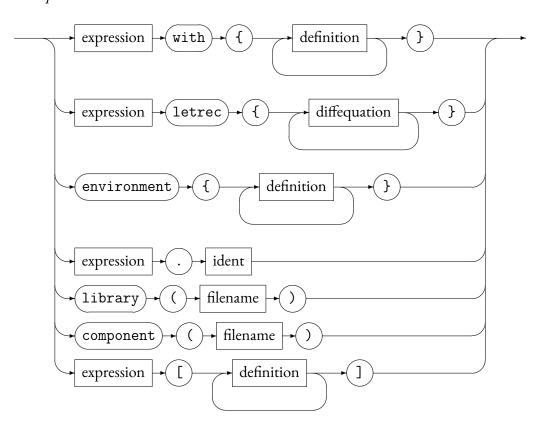
```
process = _ @ hslider("delay",0, 0, 100, 1);
```

3.4.4 Environment expressions

FAUST is a lexically scoped language. The meaning of a FAUST expression is determined by its context of definition (its lexical environment) and not by its context of use.

To keep their original meaning, FAUST expressions are bounded to their lexical environment in structures called *closures*. The following constructions allow to explicitly create and access such environments. Moreover they provide powerful means to reuse existing code and promote modular design.

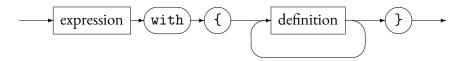
envexp



With

The with construction allows to specify a *local environment*, a private list of definition that will be used to evaluate the left hand expression

withexpression



In the following example:

```
g(x) = 1.80116083982126*x
- 0.80257737639225*x';
};
```

the definitions of f(x) and g(x) are local to $f: + \tilde{g}$.

Please note that with is left associative and has the lowest priority:

```
f: + ~ g with {...} is equivalent to (f: + ~ g) with {...}.
f: + ~ g with {...} with {...} is equivalent to ((f: + ~ g) with {...}) with {...}.
```

Letrec

The letrec construction is somehow similar to with, but for *difference equations* instead of regular definitions. It allows to easily express groups of mutually recursive signals, for example:

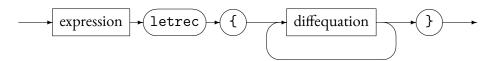
$$x(t) = y(t-1) + 10;$$

 $y(t) = x(t-1) - 1;$

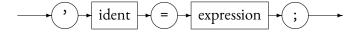
```
as E letrec { 'x = y+10; 'y = x-1; }
```

The syntax is defined by the following rules:

letrecexpression



diffequation



Please remarks the special notation 'x=y+10 instead of x=y'+10. It makes syntactically impossible to write non-sensical equations like x=x+1.

Here is a more involved example. Let say we want to define an envelop generator with an attack time, a release time and a gate signal. A possible definition is the following:

```
ar(a,r,g) = v
letrec {
   'n = (n+1) * (g<=g');
   'v = max(0, v + (n<a)/a - (n>=a)/r) * (g<=g');
};</pre>
```

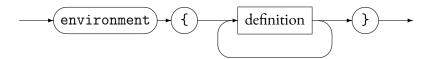
With the following semantics for n(t) and v(t):

```
\begin{array}{lll} n(t) & = & (n(t-1)+1)*(g(t) <= g(t-1)) \\ v(t) & = & \max(0,v(t-1)+(n(t-1) < a(t))/a(t) - (n(t-1) >= a(t))/r(t))*(g(t) <= g(t-1)) \\ \end{array}
```

Environment

The environment construction allows to create an explicit environment. It is like a with, but without the left hand expression. It is a convenient way to group together related definitions, to isolate groups of definitions and to create a name space hierarchy.

environment



In the following example an **environment** construction is used to group together some constant definitions:

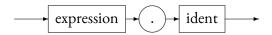
```
constant = environment {
    pi = 3.14159;
    e = 2,718;
    ...
};
```

The . construction allows to access the definitions of an environment (see next paragraph).

Access

Definitions inside an environment can be accessed using the '.' construction.

access



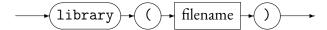
For example constant.pi refers to the definition of pi in the above constant environment

Please note that environment don't have to be named. We could have written directly environment{pi = 3.14159; e = 2,718;....}.pi

Library

The library construct allows to create an environment by reading the definitions from a file

library



For example library("miscfilter.lib") represents the environment obtained by reading the file "miscfilter.lib". It works like import("miscfilter.lib") but all the read definitions are stored in a new separate lexical environment. Individual definitions can be accessed as described in the previous paragraph. For example library ("miscfilter.lib").lowpass denotes the function lowpass as defined in the file "miscfilter.lib".

To avoid name conflicts when importing libraries it is recommended to prefer library to import. So instead of:

```
import("miscfilter.lib");
    ...
...lowpass....
...
};
```

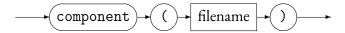
the following will ensure an absence of conflicts:

```
fl = library("miscfilter.lib");
    ...
...fl.lowpass....
...
};
```

Component

The component (...) construction allows to reuse a full FAUST program as a simple expression.

component



For example component ("freeverb.dsp") denotes the signal processor defined in file "freeverb.dsp".

Components can be used within expressions like in:

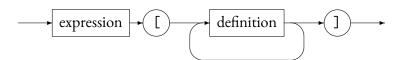
```
... component ("karplus32.dsp"): component ("freeverb.dsp")
...
```

Please note that component("freeverb.dsp") is equivalent to library("freeverb.dsp").process.

Explicit substitution

Explicit substitution can be used to customize a component or any expression with a lexical environment by replacing some of its internal definitions, without having to modify it.

explicitsubst



For example we can create a customized version of component("freeverb.dsp"), with a different definition of foo(x), by writing:

```
...component("freeverb.dsp")[foo(x) = ...;]...
};
```

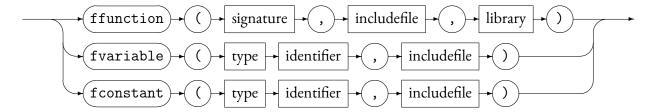
3.4.5 Foreign expressions

Reference to external C functions, variables and constants can be introduced using the foreign function mechanism.

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foreignexp



ffunction

An external C function is declared by indicating its name and signature as well as the required include file. The file "math.lib" of the FAUST distribution contains several foreign function definitions, for example the inverse hyperbolic sine function asinh:

```
asinh = ffunction(float asinh (float), <math.h>, "");
```

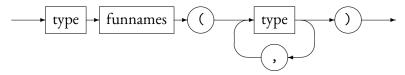
Foreign functions with input parameters are considered pure math functions. They are therefore considered free of side effects and called only when their parameters change (that is at the rate of the fastest parameter).

Exceptions are functions with no input parameters. A typical example is the C rand() function. In this case the compiler generate code to call the function at sample rate.

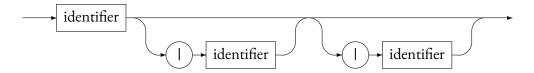
signature

The signature part (float asinh (float) in our previous example) describes the prototype of the C function: return type, function name and list of parameter types. Because the name of the foreign function can possibly depend on the floating point precision in use (float, double and quad), it is possible to give a different function name for each floating point precision using a signature with up to three function names.

signature



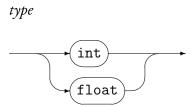
funnames



For example in the declaration asinh = ffunction(float asinhf|asinh|asinhl (float), <math.h>, "");, the signature float asinhf|asinhl (float) indicates to use the function name asinhf in single precision, asinh in double precision and asinhl in long double (quad) precision.

types

Note that currently only numerical functions involving simple int and float parameters are allowed. No vectors, tables or data structures can be passed as parameters or returned.



variables and constants

External variables and constants can also be declared with a similar syntax. In the same "math.lib" file, the definition of the sampling rate constant SR and the definition of the block-size variable BS can be found:

Foreign constants are not supposed to vary. Therefore expressions involving only foreign constants are only computed once, during the initialization period.

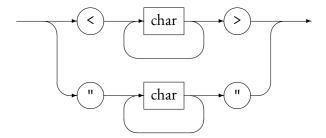
Variable are considered to vary at block speed. This means that expressions depending of external variables are computed every block.

include file

In declaring foreign functions one has also to specify the include file. It allows the FAUST compiler to add the corresponding #include... in the generated code.

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includefile



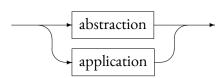
library file

In declaring foreign functions one can possibly specify the library where the actual code is located. It allows the FAUST compiler to (possibly) automatically link the library. Note that this feature is only used with the LLVM backend in 'libfaust' dynamic library model.

Applications and Abstractions

Abstractions and applications are fundamental programming constructions directly inspired by the Lambda-Calculus. These constructions provide powerful ways to describe and transform block-diagrams algorithmically.

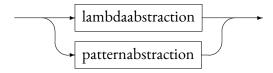




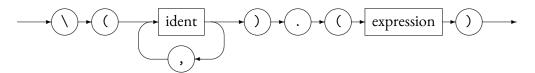
Abstractions

Abstractions correspond to functions definitions and allow to generalize a block-diagram by making variable some of its parts.

abstraction



lambdaabstraction



Let's say you want to transform a stereo reverb, freeverb for instance, into a mono effect. You can write the following expression:

```
_ <: freeverb :> _
```

The incoming mono signal is splitted to feed the two input channels of the reverb, while the two output channels of the reverb are mixed together to produce the resulting mono output.

Imagine now that you are interested in transforming other stereo effects. It can be interesting to generalize this principle by making freeverb a variable:

```
\(freeverb).(_ <: freeverb :> _)
```

The resulting abstraction can then be applied to transform other effects. Note that if freeverb is a perfectly valid variable name, a more neutral name would probably be easier to read like:

```
\(fx).(_ <: fx :> _)
```

Moreover it could be convenient to give a name to this abstraction:

```
mono = \(fx).(_ <: fx :> _);
```

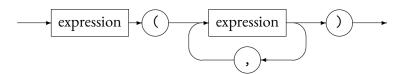
Or even use a more traditional, but equivalent, notation:

```
mono(fx) = _ <: fx :> _;
```

Applications

Applications correspond to function calls and allow to replace the variable parts of an abstraction with the specified arguments.

application



For example you can apply the previous abstraction to transform your stereo harmonizer:

```
mono(harmonizer)
```

The compiler will start by replacing mono by its definition:

```
\(fx).(_ <: fx :> _)(harmonizer)
```

Whenever the FAUST compiler find an application of an abstraction it replaces the *variable part* with the argument. The resulting expression is as expected:

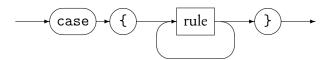
Replacing the *variable* part with the argument is called β -reduction in Lambda-Calculus

```
(_ <: harmonizer :> _)
```

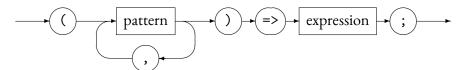
Pattern Matching

Pattern matching rules provide an effective way to analyze and transform block-diagrams algorithmically.

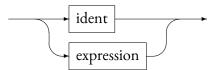
patternabstraction



Rule



Pattern



For example case { (x:y) = y:x; (x) = x; } contains two rules. The first one will match a sequential expression and invert the two part. The second one will match all remaining expressions and leave it untouched. Therefore the application:

```
case{(x:y) => y:x; (x) => x;}(freeverb:harmonizer)
```

will produce:

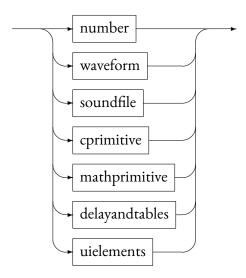
```
(harmonizer:freeverb)
```

Please note that patterns are evaluated before the pattern matching operation. Therefore only variables that appear free in the pattern are binding variables during pattern matching.

3.5 Primitives

The primitive signal processing operations represent the built-in functionalities of FAUST, that is the atomic operations on signals provided by the language. All these primitives denote *signal processors*, functions transforming *input signals* into *output signals*.

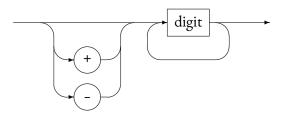
primitive



3.5.1 Numbers

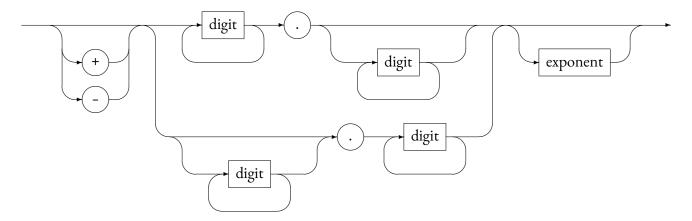
FAUST considers two types of numbers: *integers* and *floats*. Integers are implemented as 32-bits integers, and floats are implemented either with a simple, double or extended precision depending of the compiler options. Floats are available in decimal or scientific notation.

int

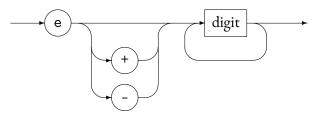


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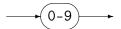
float



exponent



digit

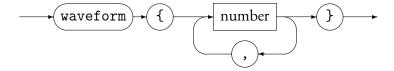


Like any other Faust expression, numbers are signal processors. For example the number 0.95 is a signal processor of type $\mathbb{S}^0 \to \mathbb{S}^1$ that transforms an empty tuple of signals () into a 1-tuple of signals (y) such that $\forall t \in \mathbb{N}, y(t) = 0.95$.

3.5.2 Waveforms

A waveform is a fixed periodic signal defined by a list of samples. A waveform has two outputs. The first output is constant and indicates the size (number of samples) of the period. The second output is the periodic signal itself.

waveform



For example waveform {0,1,2,3} produces two outputs, the constant signal 4 and the periodic signal 0,1,2,3,0,1,2,3,0,1....

Please note that waveform works nicely with rdtable. Its first output, known at compile time, gives the size of the table, while the second signal gives the content of the table. Here is an example:

```
process = waveform {10,20,30,40,50,60,70}, %(7)~+(3) :
   rdtable;
```

3.5.3 Soundfiles

The soundfile("label[url:\{pathi';'pathi';'pathi';'pathi';'pathi';'pathi';']", n)' primitive allows to access a list of externally defined sound resources, described as the list of their filename, or complete paths. The soundfile("label[url:path]", n) simplified syntax allows to use a single file.

A soundfile has:

- two inputs: the sound number (as a integer between 0 and 255 checked at compilation time), and the read index in the sound (which will access the last sample of the sound if the read index is greater than the sound length)
- two fixed outputs: the first one is the currently accessed sound length in frames, the second one is the currently accessed sound nominal sample rate in frames
- several more outputs for the sound channels themselves

If more outputs than the actual number of channels in the soundfile are used, the sound channels will be automatically duplicated up to the wanted number of outputs (so for instance if a stereo sound is used with four output channels, the same group of two channels will be duplicated).

If the soundfile cannot be loaded for whatever reason, a default sound with one channel, a length of 1024 frames and null outputs (with samples of value 0) will be used. Note also that soundfiles are entirely loaded in memory by the architecture file, so that the read index signal can access any sample.

Architecture files are responsible to load the actual soundfile. The SoundUI C++ class located in the faust/gui/SoundUI.h file implements the void addSoundfile(label, url, sf_zone) method, which loads the actual soundfiles using the *libsndfile* library, or possibly specific audio file loading code (in the case of the JUCE framework for

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instance), and set up the sf_zone sound memory pointers. If *label* is used without any *url* metadata, it will be considered as the soundfile pathname.

Note that a special architecture file can perfectly decide to access and use sound resources created by another means (that is, not directly loaded from a soundfile). For instance a mapping between labels and sound resources defined in memory could be used, with some additional code in charge to actually setup all sound memory pointers when void addSoundfile(label, url, sf_zone) is called by the buidUserInterface mechanism.

3.5.4 C-equivalent primitives

Most FAUST primitives are analogue to their C counterpart but lifted to signal processing. For example + is a function of type $\mathbb{S}^2 \to \mathbb{S}^1$ that transforms a pair of signals (x_1, x_2) into a 1-tuple of signals (y) such that $\forall t \in \mathbb{N}, y(t) = x_1(t) + x_2(t)$. The function has type $\mathbb{S}^2 \to \mathbb{S}^1$ and transforms a pair of signals (x_1, x_2) into a 1-tuple of signals (y)such that $\forall t \in \mathbb{N}, y(t) = x_1(t) - x_2(t)$.

Please be aware that the unary - only exists in a limited form. It can be used with numbers: -0.5 and variables: -myvar, but not with expressions surrounded by parenthesis, because in this case it represents a partial application. For instance - (a * b) is a partial application. It is syntactic sugar for _, (a * b): -. If you want to negate a complex operatior - only exists term in parenthesis, you'll have to use 0 - (a * b) instead.

Warning: unlike other programming languages the unary in limited form in **FAUST**

	-	
Syntax	Туре	Description
$\mid n \mid$	$\mathbb{S}^0 \to \mathbb{S}^1$	integer number: $y(t) = n$
n.m	$\mathbb{S}^0 \to \mathbb{S}^1$	floating point number: $y(t) = n.m$
_	$\mathbb{S}^1 \to \mathbb{S}^1$	identity function: $y(t) = x(t)$
!	$\mathbb{S}^1 \to \mathbb{S}^0$	cut function: $\forall x \in \mathbb{S}, (x) \to ()$
int	$\mathbb{S}^1 \to \mathbb{S}^1$	cast into an int signal: $y(t) = (int)x(t)$
float	$\mathbb{S}^1 \to \mathbb{S}^1$	cast into an float signal: $y(t) = (float)x(t)$
+	$\mathbb{S}^2 \to \mathbb{S}^1$	addition: $y(t) = x_1(t) + x_2(t)$
-	$\mathbb{S}^2 \to \mathbb{S}^1$	subtraction: $y(t) = x_1(t) - x_2(t)$
*	$\mathbb{S}^2 \to \mathbb{S}^1$	multiplication: $y(t) = x_1(t) * x_2(t)$
\wedge	$\mathbb{S}^2 \to \mathbb{S}^1$	power: $y(t) = x_1(t)^{x_2(t)}$
/	$\mathbb{S}^2 \to \mathbb{S}^1$	division: $y(t) = x_1(t)/x_2(t)$
%	$\mathbb{S}^2 \to \mathbb{S}^1$	modulo: $y(t) = x_1(t) \% x_2(t)$
&	$\mathbb{S}^2 \to \mathbb{S}^1$	logical AND: $y(t) = x_1(t) \& x_2(t)$
	$\mathbb{S}^2 \to \mathbb{S}^1$	logical OR: $y(t) = x_1(t) x_2(t)$
xor	$\mathbb{S}^2 \to \mathbb{S}^1$	logical XOR: $y(t) = x_1(t) \wedge x_2(t)$
<<	$\mathbb{S}^2 \to \mathbb{S}^1$	arith. shift left: $y(t) = x_1(t) \ll x_2(t)$
>>	$\mathbb{S}^2 \to \mathbb{S}^1$	arith. shift right: $y(t) = x_1(t) >> x_2(t)$
<	$\mathbb{S}^2 \to \mathbb{S}^1$	less than: $y(t) = x_1(t) < x_2(t)$
<=	$\mathbb{S}^2 \to \mathbb{S}^1$	less or equal: $y(t) = x_1(t) \ll x_2(t)$
>	$\mathbb{S}^2 \to \mathbb{S}^1$	greater than: $y(t) = x_1(t) > x_2(t)$
>=	$\mathbb{S}^2 \to \mathbb{S}^1$	greater or equal: $y(t) = x_1(t) >= x_2(t)$
==	$\mathbb{S}^2 \to \mathbb{S}^1$	equal: $y(t) = x_1(t) == x_2(t)$
!=	$\mathbb{S}^2 \to \mathbb{S}^1$	different: $y(t) = x_1(t)! = x_2(t)$

3.5.5 math.h-equivalent primitives

Most of the C math.h functions are also built-in as primitives (the others are defined as external functions in file math.lib).

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Syntax	Туре	Description
acos	$\mathbb{S}^1 \to \mathbb{S}^1$	arc cosine: y(t) = acosf(x(t))
asin	$\mathbb{S}^1 \to \mathbb{S}^1$	arc sine: y(t) = asinf(x(t))
atan	$\mathbb{S}^1 \to \mathbb{S}^1$	arc tangent: $y(t) = \operatorname{atanf}(x(t))$
atan2	$\mathbb{S}^2 \to \mathbb{S}^1$	arc tangent of 2 signals: $y(t) = \operatorname{atan2f}(x_1(t), x_2(t))$
cos	$\mathbb{S}^1 \to \mathbb{S}^1$	cosine: $y(t) = \cos(x(t))$
sin	$\mathbb{S}^1 \to \mathbb{S}^1$	sine: $y(t) = \sin(x(t))$
tan	$\mathbb{S}^1 \to \mathbb{S}^1$	tangent: $y(t) = \tan(x(t))$
exp	$\mathbb{S}^1 \to \mathbb{S}^1$	base-e exponential: $y(t) = \exp(x(t))$
log	$\mathbb{S}^1 \to \mathbb{S}^1$	base-e logarithm: $y(t) = \log f(x(t))$
log10	$\mathbb{S}^1 \to \mathbb{S}^1$	base-10 logarithm: $y(t) = \log 10 f(x(t))$
pow	$\mathbb{S}^2 \to \mathbb{S}^1$	power: $y(t) = powf(x_1(t), x_2(t))$
sqrt	$\mathbb{S}^1 \to \mathbb{S}^1$	square root: $y(t) = \operatorname{sqrtf}(x(t))$
abs	$\mathbb{S}^1 \to \mathbb{S}^1$	absolute value (int): $y(t) = abs(x(t))$
		absolute value (float): $y(t) = fabsf(x(t))$
min	$\mathbb{S}^2 \to \mathbb{S}^1$	$minimum: y(t) = min(x_1(t), x_2(t))$
max	$\mathbb{S}^2 \to \mathbb{S}^1$	maximum: $y(t) = \max(x_1(t), x_2(t))$
fmod	$\mathbb{S}^2 \to \mathbb{S}^1$	float modulo: $y(t) = \text{fmodf}(x_1(t), x_2(t))$
remainder	$\mathbb{S}^2 \to \mathbb{S}^1$	float remainder: $y(t) = \text{remainderf}(x_1(t), x_2(t))$
floor	$\mathbb{S}^1 \to \mathbb{S}^1$	$largest int \le: y(t) = floorf(x(t))$
ceil	$\mathbb{S}^1 \to \mathbb{S}^1$	smallest int $\geq : y(t) = \operatorname{ceilf}(x(t))$
rint	$\mathbb{S}^1 \to \mathbb{S}^1$	closest int: $y(t) = rintf(x(t))$

3.5.6 Delay, Table, Selector primitives

The following primitives allow to define fixed delays, read-only and read-write tables and 2 or 3-ways selectors (see figure 3.7).

Syntax	Type	Description
mem	$\mathbb{S}^1 \to \mathbb{S}^1$	I-sample delay: $y(t + 1) = x(t), y(0) = 0$
prefix	$\mathbb{S}^2 \to \mathbb{S}^1$	I-sample delay: $y(t+1) = x_2(t), y(0) = x_1(0)$
@	$\mathbb{S}^2 \to \mathbb{S}^1$	fixed delay: $y(t + x_2(t)) = x_1(t), y(t < x_2(t)) = 0$
rdtable	$\mathbb{S}^3 \to \mathbb{S}^1$	read-only table: $y(t) = T[r(t)]$
rwtable	$\mathbb{S}^5 \to \mathbb{S}^1$	read-write table: $T[w(t)] = c(t); y(t) = T[r(t)]$
select2	$\mathbb{S}^3 \to \mathbb{S}^1$	select between 2 signals: $T[] = \{x_0(t), x_1(t)\}; y(t) = T[s(t)]$
select3	$\mathbb{S}^4 \to \mathbb{S}^1$	select between 3 signals: $T[]=\{x_0(t),x_1(t),x_2(t)\};y(t)=T[s(t)]$

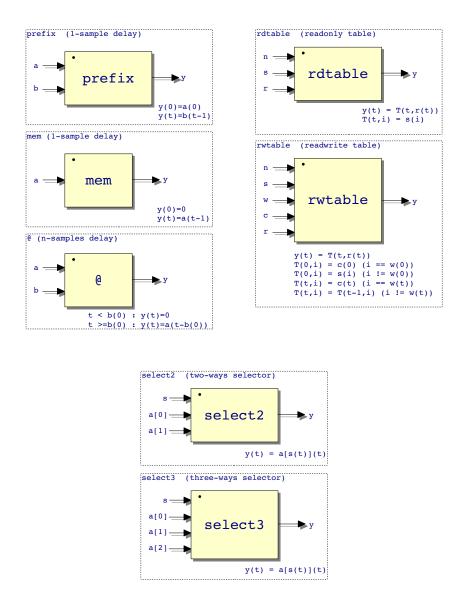


Figure 3.7: Delays, tables and selectors primitives

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3.5.7 User Interface Elements

FAUST user interface widgets allow an abstract description of the user interface from within the FAUST code. This description is independent of any GUI toolkits. It is based on *buttons*, *checkboxes*, *sliders*, etc. that are grouped together vertically and horizontally using appropriate grouping schemes.

All these GUI elements produce signals. A button for example (see figure 3.8) produces a signal which is 1 when the button is pressed and 0 otherwise. These signals can be freely combined with other audio signals.

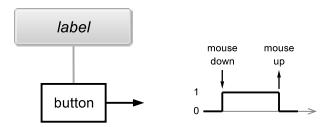


Figure 3.8: User Interface Button

Syntax	Example
button(str)	button("play")
checkbox(str)	checkbox("mute")
<pre>vslider(str, cur, min, max, step)</pre>	vslider("vol",50,0,100,1)
hslider(str, cur, min, max, step)	hslider("vol",0.5,0,1,0.01)
<pre>nentry(str, cur, min, max, step)</pre>	nentry("freq",440,0,8000,1)
vgroup(str, block-diagram)	vgroup("reverb",)
hgroup(str, block-diagram)	hgroup("mixer",)
tgroup(str, block-diagram)	tgroup("parametric",)
vbargraph(str, min, max)	vbargraph("input",0,100)
hbargraph(str, min, max)	hbargraph("signal",0,1.0)
attach	attach(x, vumeter(x))

All numerical parameters (like *cur*, *min*, *max*, *step*) have to be known at compile time.

Labels

Every user interface widget has a label (a string) that identifies it and informs the user of its purpose. There are three important mechanisms associated with labels (and coded inside the string): *variable parts*, *pathnames* and *metadata*.

Variable parts. Labels can contain variable parts. These variable parts are indicated by the sign "%" followed by the name of a variable. During compilation each label is processed in order to replace the variable parts by the value of the variable. For example par(i,8,hslider("Voice %i", 0.9, 0, 1, 0.01)) creates 8 different sliders in parallel.

```
hslider("Voice 0", 0.9, 0, 1, 0.01),
hslider("Voice 1", 0.9, 0, 1, 0.01),
...
hslider("Voice 7", 0.9, 0, 1, 0.01).
```

while par(i,8,hslider("Voice", 0.9, 0, 1, 0.01)) would have created only one slider and duplicated its output 8 times.

The variable part can have an optional format digit. For example "Voice %2i" would indicate to use two digit when inserting the value of i in the string.

An escape mechanism is provided. If the sign % is followed by itself, it will be included in the resulting string. For example "feedback (%%)" will result in "feedback (%)".

Pathnames. Thanks to horizontal, vertical and tabs groups, user interfaces have a hierarchical structure analog to a hierarchical file system. Each widget has an associated *pathname* obtained by concatenating the labels of all its surrounding groups with its own label.

In the following example:

```
hgroup("Foo",
...
vgroup("Faa",
...
hslider("volume",...)
...
)
...
```

the volume slider has pathname /h:Foo/v:Faa/volume.

In order to give more flexibility to the design of user interfaces, it is possible to explicitly specify the absolute or relative pathname of a widget directly in its label.

In our previous example the pathname of :

```
hslider("../volume",...)
```

would have been "/h:Foo/volume", while the pathname of:

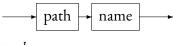
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```
hslider("t:Fii/volume",...)
```

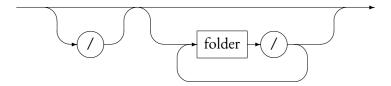
would have been: "/h:Foo/v:Faa/t:Fii/volume".

The grammar for labels with pathnames is the following:

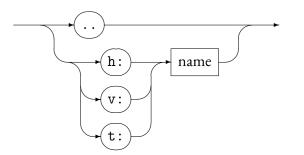
label



path



folder



Metadata Widget labels can contain metadata enclosed in square brackets. These metadata associate a key with a value and are used to provide additional information to the architecture file. They are typically used to improve the look and feel of the user interface. The FAUST code:

will produce and the corresponding C++ code:

```
class mydsp : public dsp {
    ...
    virtual void buildUserInterface(UI* interface) {
        interface->openVerticalBox("m");
        interface->declare(&fslider0, "key1", "val 1");
        interface->declare(&fslider0, "key2", "val 2");
}
```

```
interface -> addHorizontalSlider("foo",
    &fslider0, 0.0f, 0.0f, 1.0f, 0.1f);
   interface -> closeBox();
}
...
};
```

All the metadata are removed from the label by the compiler and transformed in calls to the UI::declare() method. All these UI::declare() calls will always take place before the UI::AddSomething() call that creates the User Interface element. This allows the UI::AddSomething() method to make full use of the available metadata.

It is the role of the architecture file to decide what to do with these metadata. The <code>jack-qt.cpp</code> architecture file for example implements the following:

- I. "...[style:knob]..." creates a rotating knob instead of a regular slider or nentry.
- 2. "...[style:led]..." in a bargraph's label creates a small LED instead of a full bargraph
- 3. "...[unit:dB]..." in a bargraph's label creates a more realistic bargraph with colors ranging from green to red depending of the level of the value
- 4. "... [unit:xx]..." in a widget postfixes the value displayed with xx
- 5. "... [tooltip:bla bla]..." add a tooltip to the widget
- 6. "...[osc:/address min max]..." Open Sound Control message alias

Moreover starting a label with a number option like in "[1]..." provides a convenient means to control the alphabetical order of the widgets.

Attach

The attach primitive takes two input signals and produce one output signal which is a copy of the first input. The role of attach is to force its second input signal to be compiled with the first one. From a mathematical point of view attach(x,y) is equivalent to 1*x+0*y, which is in turn equivalent to x, but it tells the compiler not to optimize-out y.

To illustrate this role let say that we want to develop a mixer application with a vumeter for each input signals. Such vumeters can be easily coded in FAUST using an envelop detector connected to a bargraph. The problem is that these envelop signals have no role in the output signals. Using attach(x,vumeter(x)) one can tell the compiler that when x is compiled vumeter(x) should also be compiled.

Chapter 4

Invoking the Faust compiler

The Faust compiler is invoked using the faust command. It translate Faust programs into C++ code. The generated code can be wrapped into an optional *architecture file* allowing to directly produce a fully operational program.

compiler



For example faust noise.dsp will compile noise.dsp and output the corresponding C++ code on the standard output. The option -o allows to choose the output file: faust noise.dsp -o noise.cpp. The option -a allows to choose the architecture file: faust -a alsa-gtk.cpp noise.dsp.

To compile a FAUST program into an ALSA application on Linux you can use the following commands:

```
faust -a alsa-gtk.cpp noise.dsp -o noise.cpp
g++ -lpthread -lasound
    'pkg-config --cflags --libs gtk+-2.0'
    noise.cpp -o noise
```

4.1 Structure of the generated code

A FAUST DSP C++ class derives from the base *dsp* class defined as below:

```
class dsp {
    public:
        dsp() {}
        virtual ~dsp() {}
        virtual int getNumInputs() = 0;
        virtual int getNumOutputs() = 0
        virtual void buildUserInterface(UI* ui_interface) =
           0:
        virtual int getSampleRate() = 0;
        virtual void init(int samplingRate) = 0;
        virtual void instanceInit(int samplingRate) = 0;
        virtual void instanceConstants(int samplingRate) = 0;
        virtual void instanceResetUserInterface() = 0;
        virtual void instanceClear() = 0;
        virtual dsp* clone() = 0;
        virtual void metadata(Meta* m) = 0;
        virtual void compute(int count, FAUSTFLOAT** inputs,
           FAUSTFLOAT** outputs) = 0;
        virtual void compute(double date_usec, int count,
           FAUSTFLOAT** inputs, FAUSTFLOAT** outputs)
        {
            compute(count, inputs, outputs);
        }
};
```

Here is the class generated with the command faust noise.dsp. Methods are filled by the compiler with the actual code.

Several fine-grained initialization methods are available. The instanceInit method calls several additional initialization methods. The instanceConstants method sets the instance constant state. The instanceClear method resets the instance dynamic state (delay lines...). The instanceResetUserInterface method resets all control value to their default state. All of those methods can be used individually on an allocated instance to reset part of its state.

The classInit static method will initialize static tables that are shared between all instances of the class, and is typically supposed to be called once.

Finally the init method combines class static state and instance initialization.

When using a single instance, then calling init is the simplest way to do what is needed.

When using several instances, then all of them can be initialized using instanceInit, whith a single call to classInit to initialize the static shared state.

```
class mydsp : public dsp {
  private:
   FAUSTFLOAT fslider0;
         iRec0[2];
    int
    int fSamplingFreq;
  public:
    virtual void metadata(Meta* m) {
       m->declare("name", "Noise");
        m->declare("version", "1.1");
       m->declare("author", "Grame");
       m->declare("license", "BSD");
       m->declare("copyright", "(c)GRAME 2009");
    }
    virtual int getNumInputs() { return 0; }
    virtual int getNumOutputs() { return 1; }
    static void classInit(int samplingFreq) {
    virtual void instanceConstants(int samplingFreq) {
        fSamplingFreq = samplingFreq;
    virtual void instanceResetUserInterface() {
        fslider0 = 0.5f;
    virtual void instanceClear() {
       for (int i=0; i<2; i++) iRec0[i] = 0;
    virtual void init(int samplingFreq) {
        classInit(samplingFreq);
        instanceInit(samplingFreq);
    virtual void instanceInit(int samplingFreq) {
        instanceConstants(samplingFreq);
        instanceResetUserInterface();
        instanceClear();
    virtual mydsp* clone() {
        return new mydsp();
    virtual int getSampleRate() {
        return fSamplingFreq;
    virtual void buildUserInterface(UI* ui_interface) {
        ui_interface -> openVerticalBox("Noise");
```

4.2 Compilation options

Compilation options are listed in the following table:

Short	Long	Description
-h	-help	print the help message
-v	-version	print version information
-d	-details	print compilation details
-tg	-task-graph	draw a graph of all internal compu-
		tation loops as a .dot (graphviz) file.
-sg	-signal-graph	draw a graph of all internal signal
		expressions as a .dot (graphviz) file.
-ps	-postscript	generate block-diagram postscript
		files
-svg	-svg	generate block-diagram svg files
-blur	-shadow-blur	add a blur to boxes shadows
-sd	-simplify-diagrams	simplify block-diagram before
		drawing them
-f n	-fold n	max complexity of svg diagrams be-
		fore splitting into several files (de-
		fault 25 boxes)
		continued on next page

-mns n -max-name-size n max character size used in syg diagram labels use simple names (without arguments) for block-diagram (default max size: 40 chars) generate an additional description file in xml format -uim -user-interface-macros add user interface macro definitions to the C++ code list all the source files and libraries implied in a compilation prints the internal signals in normalized-form prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -rb -right-balanced generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default 32 samples) when-vec -vec -vectorize generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using OpenMP (implies -vec) generate parallel code using threads directly (implies -vec) generate parallel code using threads directly (implies -vec) group sequential tasks together when-omp or-sch is used -single -single-precision-floatuse doubles for internal computations (default)	Short	Long	Description
-sm -simple-names use simple names (without arguments) for block-diagram (default max size: 40 chars) generate an additional description file in xml format -uim -user-interface-macros add user interface macro definitions to the C++ code -flist -file-list list all the source files and libraries implied in a compilation prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -rb -less-temporaries generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default), 1:simple] when -vec loop variant [0:fastest (default), 1:simple] when -vec schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatme doubles for internal computations (default) -double -double-precision-floatme doubles for internal computations	-mns n	-max-name-size n	max character size used in svg dia-
ments) for block-diagram (default max size: 40 chars) generate an additional description file in xml format -uim -user-interface-macros add user interface macro definitions to the C++ code -flist -file-list list all the source files and libraries implied in a compilation -norm -normalized-form prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions -mb -mid-balanced generate left-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -ncd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code -vs n -vec-size n size of the vector (default 32 samples) when -vec loop variant [o:fastest (default), 1:simple] when -vec -dfs -deepFirstScheduling generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatme doubles for internal computations (default) -double -double-precision-floatme doubles for internal computations			gram labels
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-xml -xml generate an additional description file in xml format -uim -user-interface-macros add user interface macro definitions to the C++ code -flist -file-list list all the source files and libraries implied in a compilation -norm -normalized-form prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions -mb -mid-balanced generate mid-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -less-temporaries generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), r.simple] when -vec -dfs -deepFirstScheduling Generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations			ments) for block-diagram (default
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-uim -user-interface-macros add user interface macro definitions to the C++ code list all the source files and libraries implied in a compilation prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions generate mid-balanced expressions (default) -rb -right-balanced generate right-balanced expressions generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), r:simple] when -vec -dfs -deepFirstScheduling openMP (implies -vec) -sch -scheduler generate parallel code using OpenMP (implies -vec) generate parallel code using threads directly (implies -vec) generate parallel code using threads directly (implies -vec) group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations	-xml	-xml	generate an additional description
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-flist -file-list list all the source files and libraries implied in a compilation prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -rb -right-balanced generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), 1:simple] when -vec schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse doubles for internal computations (default) -double -double-precision-floatuse doubles for internal computations	-uim	-user-interface-macro	add user interface macro defini-
implied in a compilation prints the internal signals in normalized form and exits -lb			tions to the C++ code
-norm -normalized-form prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -rb -less-temporaries generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code -vs n -vec-size n size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), 1:simple] when -vec -dfs -deepFirstScheduling openAP (implies -vec) -sch -scheduler generate parallel code using OpenMP (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations	-flist	-file-list	list all the source files and libraries
-norm -normalized-form prints the internal signals in normalized form and exits -lb -left-balanced generate left-balanced expressions (default) -rb -right-balanced generate right-balanced expressions (default) -rb -less-temporaries generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code -vs n -vec-size n size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), 1:simple] when -vec -dfs -deepFirstScheduling openAP (implies -vec) -sch -scheduler generate parallel code using OpenMP (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations			implied in a compilation
malized form and exits -lb	-norm	-normalized-form	
-mb -mid-balanced generate mid-balanced expressions (default) -rb -right-balanced generate right-balanced expressions generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), r:simple] when -vec -dfs -deepFirstScheduling schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations			^
-mb -mid-balanced generate mid-balanced expressions (default) -rb -right-balanced generate right-balanced expressions generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), r:simple] when -vec -dfs -deepFirstScheduling schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations	-1b	-left-balanced	generate left-balanced expressions
-rb -right-balanced generate right-balanced expressions generate less temporaries in compiling delays -mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code -vs n -vec-size n size of the vector (default 32 samples) when -vec loop variant [o:fastest (default), r:simple] when -vec schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatse doubles for internal computations (default) -double -double-precision-floatse doubles for internal computations	-mb	-mid-balanced	1 0
-less-temporaries generate less temporaries in compiling delays threshold between copy and ring buffer delays (default 16 samples) -vec			1 0 1
-less-temporaries generate less temporaries in compiling delays threshold between copy and ring buffer delays (default 16 samples) -vec	-rb	-right-balanced	generate right-balanced expressions
ing delays threshold between copy and ring buffer delays (default 16 samples) -vec -vec-size n size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), 1:simple] when -vec schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) generate parallel code using threads directly (implies -vec) group sequential tasks together when -omp or -sch is used -single -single -precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computa- tions	-lt	_	
-mcd n -max-copy-delay n threshold between copy and ring buffer delays (default 16 samples) -vec -vectorize generate easier to vectorize code -vs n -vec-size n size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), 1:simple] when -vec -dfs -deepFirstScheduling schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatuse floats for internal computations (default) -double -double-precision-floatuse doubles for internal computations		-	
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-vec -vectorize generate easier to vectorize code -vs n -vec-size n size of the vector (default 32 samples) when -vec -lv n -loop-variant n loop variant [o:fastest (default), 1:simple] when -vec -dfs -deepFirstScheduling schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatse floats for internal computations (default) -double -double-precision-floatse doubles for internal computations			
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-dfs -deepFirstScheduling schedule vector loops in deep first order when -vec -omp -openMP generate parallel code using OpenMP (implies -vec) -sch -scheduler generate parallel code using threads directly (implies -vec) -g -groupTasks group sequential tasks together when -omp or -sch is used -single -single-precision-floatse floats for internal computations (default) -double -double-precision-floatse doubles for internal computations		1	
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-double -double-precision-floatuse doubles for internal computations	-single	-single-precision-floa	
-double -double-precision-floatuse doubles for internal computations	0		_
tions	-double	-double-precision-floa	'
		1 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2 2	_
CONLLINACIA ON MEXIL DAME			continued on next page

Short	Long	Description
-quad	-quad-precision-float:	s use extended for internal computa-
		tions
-mdoc	-mathdoc	generates the full mathematical de-
		scription of a FAUST program
-mdlang	-mathdoc-lang $\it l$	choose the language of the mathe-
1		matical description ($l = en, fr,$)
-stripmdo	c-strip-mdoc-tags	remove documentation tags when
		printing FAUST listings
-cn	-class-name name	name of the dsp class to be used in-
name		stead of 'mydsp'
-t time	-timeout time	time out of time seconds (default
		600) for the compiler to abort
-a file		architecture file to use
-a file -o file		C++ output file

Chapter 5

Embedding the Faust compiler using libfaust

The dynamic compilation chain allows developers to embed the Faust compiler technology directly in their application or plugins. Thanks to the awesome LLVM technology combined with libfaust, the library version of the Faust compiler, Faust DSP programs can directly be compiled and executed on the fly at full speed.

5.1 Dynamic compilation chain

The FAUST compiler uses an intermediate FIR representation (FAUST Imperative Representation), which can be translated to several output languages. The FIR language describes the computation performed on the samples in a generic manner. It contains primitives to read and write variables and arrays, do arithmetic operations, and define the necessary control structures (for and while loops, if structure etc.).

To generate various output languages, several backends have been developed: for C, C++, Java, JavaScript, asm.js, and LLVM IR. The native LLVM based compilation chain is particularly interesting: it provides direct compilation of a DSP source into executable code in memory, bypassing the external compiler requirement.

5.2 LLVM

LLVM (formerly Low Level Virtual Machine) is a compiler infrastructure, designed for compile-time, link-time, run-time optimization of programs written in arbitrary programming languages. Executable code is produced dynamically using a *Just In Time*

compiler from a specific code representation, called LLVM IR. Clang, the LLVM native C/C++/Objective- C compiler is a front-end for LLVM Compiler. It can, for instance, convert a C or C++ source file into LLVM IR code. Domain-specific languages like FAUST can easily target the LLVM IR. This has been done by developing an LLVM IR backend in the FAUST compiler.

5.3 Compiling in memory

The complete chain goes from the FAUST DSP source code, compiled in LLVM IR using the LLVM backend, to finally produce the executable code using the LLVM JIT. All steps take place in memory, getting rid of the classical file based approaches. Pointers to executable functions can be retrieved from the resulting LLVM module and the code directly called with the appropriate parameters.

The FAUST compiler has been packaged as an embeddable library called libfaust, published with an associated API that imitates the concept of oriented-object languages, like C++. Given a FAUST source code (as a file or a string), calling the createDSPFactoryXXX function runs the compilation chain (FAUST + LLVM JIT) and generates the *prototype* of the class, as a llvm_dsp_factory pointer.

```
class llvm_dsp_factory {

public:

    /* Return Factory name */
    std::string getName();

    /* Return Factory LLVM target */
    std::string getTarget();

    /* Return Factory SHA key */
    std::string getSHAKey();

    /* Return Factory expanded DSP code */
    std::string getDSPCode();

    /* Create a new DSP instance, to be deleted with C++
        'delete' */
    llvm_dsp* createDSPInstance();

    /* Set a custom memory manager to be used when
        creating instances */
```

```
void setMemoryManager(dsp_memory_manager* manager);

/* Return the currently set custom memory manager */
dsp_memory_manager* getMemoryManager();
};
```

Note that the library keeps an internal cache of all allocated factories so that the compilation of the same DSP code, that is same source code and same set of *normalized* (= sorted in a canonical order) compilations options, will return the same (reference counted) factory pointer. You will have to explicitly use deleteDSPFactory to properly decrement the reference counter when the factory is no more needed. You can get a unique SHAI key of the created factory using its getSHAKey method.

Next, the createDSPInstance function, corresponding to the new className of C++, instantiates a llvm_dsp pointer to be used through its interface, connected to the audio chain and controller interfaces. When finished, use delete to destroy the dsp instance.

```
class llvm_dsp : public dsp {
 public:
    int getNumInputs();
    int getNumOutputs();
    void buildUserInterface(UI* ui_interface);
    int getSampleRate();
    void init(int samplingRate);
    void instanceInit(int samplingRate);
    void instanceConstants(int samplingRate);
    void instanceResetUserInterface();
    void instanceClear();
    llvm_dsp* clone();
    void metadata(Meta* m);
    void compute(int count, FAUSTFLOAT** inputs,
       FAUSTFLOAT** outputs);
};
```

Since 11vm_dsp is a subclass of the dsp base class, an object of this type can be used with all already available audio and UI classes, in essence reusing all architecture files already

developed for the static C++ class compilation scheme (like OSCUI, httpdUI interfaces etc.), look at Developing a new architecture file section.

5.4 Saving/restoring the factory

After the DSP factory has been compiled, your application or plugin may want to save/restore it in order to save Faust to LLVM IR compilation or even JIT compilation time at next use. To get the internal factory compiled code, several functions are available:

- writeDSPFactoryToIR allows to get the DSP factory LLVM IR (in textual format) as a string, writeDSPFactoryToIRFile allows to save the DSP factory LLVM IR (in textual format) in a file,
- writeDSPFactoryToBitcode allows to get the DSP factory LLVM IR (in binary format) as a string, writeDSPFactoryToBitcodeFile allows to save the DSP factory LLVM IR (in binary format) in a file,
- writeDSPFactoryToMachine allows to get the DSP factory executable machine code as a string, writeDSPFactoryToMachineFile allows to save the DSP factory executable machine code in a file.

To re-create a DSP factory from a previously saved code, several functions are available:

- readDSPFactoryFromIR allows to create a DSP factory from a string containing the LLVM IR (in textual format), readDSPFactoryFromIRFile allows to create a DSP factory from a file containing the LLVM IR (in textual format),
- readDSPFactoryFromBitcode allows to create a DSP factory from a string containing the LLVM IR (in binary format), readDSPFactoryFromBitcodeFile allows to create a DSP factory from a file containing the LLVM IR (in binary format),
- readDSPFactoryFromMachine allows to create a DSP factory from a string containing the executable machine code, readDSPFactoryFromMachineFile allows to create a DSP factory from a file containing the executable machine code.

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5.5 Additional functions

Some additional functions are available in the libfaust API:

expandDSPFromString/expandDSPFromFile creates a self-contained DSP source string where all needed librairies have been included. All compilations options are normalized and included as a comment in the expanded string, generateAuxFilesFromString/generateAuxFilesFromFile: from a DSP source string or file, generates auxiliary files: SVG, XML, ps... depending of the argy parameters.

5.6 Using the libfaust library

The libfaust library is part of the FAUST tree. You'll have to compile and install it. Then look at the installed faust/dsp/llvm-dsp.h header for a complete description of the API. Note that faust/dsp/llvm-c-dsp.h is a pure C version of the same API. The additional functions are available in the faust/dsp/libfaust.h header and their C version is in faust/dsp/libfaust-c.h.

5.7 Use case examples

The dynamic compilation chain has been used in several projects:

- FaustLive, an integrated IDE for Faust development
- Faustgen, an external object for Cycling Max/MSP language
- Csound6, see this demo video
- LibAudioStream, a framework to manipulate audio ressources through the concept of streams
- Oliver Larkin JUCE framework integration and pMix2 project
- · an experimental version of Antescofo
- FaucK: the combination of the Chuck language and FAUST

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Chapter 6

Architecture files

A FAUST program describes a *signal processor*, a pure computation that maps *input signals* to *output signals*. It says nothing about audio drivers or GUI toolkits. This missing information is provided by *architecture files*.

An *architecture file* describes how to relate a FAUST program to the external world, in particular the audio drivers and the user interface to be used. This approach allows a single FAUST program to be easily deployed to a large variety of audio standards (Max/MSP externals, PD externals, VST plugins, CoreAudio applications, Jack applications, iPhone, etc.).

The architecture to be used is specified at compile time with the -a options. For example faust -a jack-gtk.cpp foo.dsp indicates to use the Jack GTK architecture when compiling foo.dsp.

The main available architecture files are listed table 6.1. Since FAUST 0.9.40 some of these architectures are a modular combination of an *audio module* and one or more *user interface modules*. Among these user interface modules OSCUI provide supports for Open Sound Control allowing FAUST programs to be controlled by OSC messages.

6.1 Audio architecture modules

An *audio architecture module* typically connects a FAUST program to the audio drivers. It is responsible for allocating and releasing the audio channels and for calling the FAUST dsp::compute method to handle incoming audio buffers and/or to produce audio output. It is also responsible for presenting the audio as non-interleaved float data, normalized between -1.0 and 1.0.

A FAUST audio architecture module derives an audio class defined as below:

File name	Description
alchemy-as.cpp	Flash - ActionScript plugin
ca-qt.cpp	CoreAudio QT4 standalone application
jack-gtk.cpp	JACK GTK standalone application
jack-qt.cpp	JACK QT ₄ standalone application
jack-console.cpp	JACK command line application
jack-internal.cpp	JACK server plugin
alsa-gtk.cpp	ALSA GTK standalone application
alsa-qt.cpp	ALSA QT4 standalone application
oss-gtk.cpp	OSS GTK standalone application
pa-gtk.cpp	PortAudio GTK standalone application
pa-qt.cpp	PortAudio QT4 standalone application
max-msp.cpp	Max/MSP external
vst.cpp	VST plugin
vst2p4.cpp	VST 2.4 plugin
vsti-mono.cpp	VSTi mono instrument
vsti-poly.cpp	VSTi polyphonic instrument
ladspa.cpp	LADSPA plugin
q.cpp	Q language plugin
supercollider.cpp	SuperCollider Unit Generator
snd-rt-gtk.cpp	Snd-RT music programming language
csound.cpp	CSOUND opcode
puredata.cpp	PD external
sndfile.cpp	sound file transformation command
bench.cpp	speed benchmark
octave.cpp	Octave plugin
plot.cpp	Command line application
sndfile.cpp	Command line application

Table 6.1: Some of the available architectures.

The API is simple enough to give a great flexibility to audio architectures implementations. The init method should initialize the audio. At init exit, the system should be in a safe state to recall the dsp object state.

Table 6.2 gives some of the audio architectures currently available for various operating systems.

Audio system	Operating system
Alsa	Linux
CoreAudio	Mac OS X, iOS
JACK	Linux, Mac OS X, Windows
PortAudio	Linux, Mac OS X, Windows
OSC	Linux, Mac OS X, Windows
VST	Mac OS X, Windows
Max/MSP	Mac OS X, Windows
Csound	Linux, Mac OS X, Windows
SuperCollider	Linux, Mac OS X, Windows
PureData	Linux, Mac OS X, Windows
Pure	Linux, Mac OS X, Windows

Table 6.2: Some of FAUST audio architectures.

6.2 UI architecture modules

A UI architecture module links user actions (via graphic widgets, command line parameters, OSC messages, etc.) with the FAUST program to control. It is responsible for asso-

ciating program parameters to user interface elements and to update parameter's values according to user actions. This association is triggered by the dsp::buildUserInterface call, where the dsp asks a UI object to build the DSP module controllers.

Since the interface is basically graphic oriented, the main concepts are *widget* based: a UI architecture module is semantically oriented to handle active widgets, passive widgets and widgets layout.

A FAUST UI architecture module derives an *UI* class (Figure 6.1).

6.2.1 Active widgets

Active widgets are graphical elements that control a parameter value. They are initialized with the widget name and a pointer to the linked value, using the FAUST-FLOAT macro type (defined at compile time as either float or double). The widget currently considered are Button, CheckButton, VerticalSlider, HorizontalSlider and NumEntry.

A GUI architecture must implement a method addXxx(const char* name, FAUSTFLOAT* zone, ...) for each active widget. Additional parameters are available for Slider and NumEntry: the init, min, max and step values.

6.2.2 Passive widgets

Passive widgets are graphical elements that reflect values. Similarly to active widgets, they are initialized with the widget name and a pointer to the linked value. The widget currently considered are HorizontalBarGraph and VerticalBarGraph.

A UI architecture must implement a method

addXxx(const char* name, FAUSTFLOAT* zone, ...) for each passive widget. Additional parameters are available, depending on the passive widget type.

6.2.3 Widgets layout

Generally, a GUI is hierarchically organized into boxes and/or tab boxes. A UI architecture must support the following methods to setup this hierarchy:

```
openTabBox(const char* label)
openHorizontalBox(const char* label)
openVerticalBox(const char* label)
closeBox(const char* label)
```

Note that all the widgets are added to the current box.

```
#ifndef FAUSTFLOAT
#define FAUSTFLOAT float
#endif
class UI
 public:
          UI() {}
  virtual ~UI() {}
     -- widget layouts
  virtual void openTabBox(const char* 1) = 0;
  virtual void openHorizontalBox(const char* 1) = 0;
  virtual void openVerticalBox(const char* 1) = 0;
  virtual void closeBox() = 0;
   -- active widgets
  virtual void addButton(const char* 1, FAUSTFLOAT* z)
  virtual void addCheckButton(const char* 1, FAUSTFLOAT* z)
    = 0;
  virtual void addVerticalSlider(const char* 1,
          FAUSTFLOAT* z,
          FAUSTFLOAT init, FAUSTFLOAT min,
          FAUSTFLOAT max, FAUSTFLOAT step) = 0;
  virtual void addHorizontalSlider(const char* 1,
          FAUSTFLOAT* z,
          FAUSTFLOAT init, FAUSTFLOAT min,
          FAUSTFLOAT max, FAUSTFLOAT step) = 0;
  virtual void addNumEntry(const char* 1, FAUSTFLOAT* z,
          FAUSTFLOAT init, FAUSTFLOAT min,
          FAUSTFLOAT max, FAUSTFLOAT step) = 0;
   -- passive widgets
  virtual void addHorizontalBargraph(const char* 1,
                        FAUSTFLOAT* z, FAUSTFLOAT min,
                        FAUSTFLOAT max) = 0;
  virtual void addVerticalBargraph(const char* 1,
                        FAUSTFLOAT* z, FAUSTFLOAT min,
                        FAUSTFLOAT max) = 0;
  -- metadata declarations
  virtual void declare(FAUSTFLOAT*, const char*, const char*)
     {}
};
```

Figure 6.1: UI, the root user interface class.

UI	Comment
console	a textual command line UI
GTKUI	a GTK-based GUI
QTGUI	a multi-platform QT-based GUI
FUI	a file-based UI to store and recall modules states
OSCUI	OSC control (see section 7)
httpdUI	HTTP control (see section 8)

Table 6.3: Some of the available UI architectures.

6.2.4 Metadata

The FAUST language allows widget labels to contain metadata enclosed in square brackets as key/value pairs. These metadata are handled at GUI level by a declare method taking as argument, a pointer to the widget associated zone, the metadata key and value:

declare(FAUSTFLOAT* zone, const char* key, const char* value)

Here is the table of currently supported general medatada (look at section 7 for OSC specific metadata and section 9 for MIDI specific metadata):

Key	Value
tooltip	actual string content
hidden	0 or 1
size	actual value
unit	Hz or dB
scale	log or exp
style	knob or led or numerical
style	<pre>radio{'label1':v1;'label2':v2}</pre>
style	menu{'label1':v1;'label2':v2}
acc	axe curve amin amid amax
gyr	axe curve amin amid amax
screencolor	red or green or blue or white

Table 6.4: Supported medatada.

Some typical example where several metadata are defined could be:

```
nentry("freq [unit:Hz][scale:log][acc:0 0 -30 0 30][style:menu{'white
noise':0;'pink noise':1;'sine':2}][hidden:0]", 0, 20, 100, 1)
```

or:

```
vslider("freq [unit:dB][style:knob][gyr:0 0 -30 0 30]", 0, 20, 100,
1)
```

Note that medatada are not supported in all architecture files. Some of them like (acc or gyr for exemple) only make sense on platforms with accelerometers or gyroscopes sensors. The set of medatada may be extended in the future.

6.3 Developing a new architecture file

Developing a new architecture file typically means writing a generic C++ file, that will be populated with the actual output of the FAUST compiler, in order to produce a complete C++ file, ready to be compiled as a standalone application or plugin.

The architecture to be used is specified at compile time with the -a option. It must contain the <<includeIntrinsic>> and <<includeclass>> lines that will be looked at by the FAUST compiler, and replaced by the generated C++ class.

Look at the minimal.cpp example located in the architecture folder:

```
#include "faust/gui/PrintUI.h"
#include "faust/gui/meta.h"
#include "faust/audio/dummy-audio.h"
using std::max;
using std::min;
// FAUST generated signal processor
<<includeIntrinsic>>
<<includeclass>>
int main(int argc, char *argv[])
    mydsp DSP;
    PrintUI ui;
    // Activate the UI
    // (here that only prints the control paths)
    DSP.buildUserInterface(&ui);
    // Allocate the audio driver to render
    // 5 buffers of 512 frames
    dummyaudio audio(5);
```

```
audio.init("Test", &DSP);

// Render buffers...
audio.start();

audio.stop();
}
```

Calling faust -a minimal.cpp noise.dsp -a noise.cpp will produce a ready to compile noise.cpp file:

```
#include "faust/gui/PrintUI.h"
#include "faust/gui/meta.h"
#include "faust/audio/dummy-audio.h"
using std::max;
using std::min;
// FAUST generated signal processor
#ifndef FAUSTFLOAT
#define FAUSTFLOAT float
#endif
#ifndef FAUSTCLASS
#define FAUSTCLASS mydsp
#endif
class mydsp : public dsp {
  private:
    FAUSTFLOAT fslider0;
    int iRec0[2];
    int fSamplingFreq;
  public:
    virtual void metadata(Meta* m) {
        m->declare("name", "Noise");
        m->declare("version", "1.1");
        m->declare("author", "Grame");
m->declare("license", "BSD");
        m->declare("copyright", "(c)GRAME 2009");
    }
    virtual int getNumInputs() { return 0; }
    virtual int getNumOutputs() { return 1; }
    static void classInit(int samplingFreq) {
```

```
virtual void instanceConstants(int samplingFreq) {
        fSamplingFreq = samplingFreq;
    virtual void instanceResetUserInterface() {
        fslider0 = 0.5f;
    virtual void instanceClear() {
        for (int i=0; i<2; i++) iRec0[i] = 0;
    virtual void init(int samplingFreq) {
        classInit(samplingFreq);
        instanceInit(samplingFreq);
    }
    virtual void instanceInit(int samplingFreq) {
        instanceConstants(samplingFreq);
        instanceResetUserInterface();
        instanceClear();
    }
    virtual mydsp* clone() {
        return new mydsp();
    virtual int getSampleRate() {
        return fSamplingFreq;
    virtual void buildUserInterface(UI* ui_interface) {
        ui_interface -> openVerticalBox("Noise");
        ui_interface -> declare (&fslider0, "acc", "0 0 -10 0 10
           ");
        ui_interface ->declare(&fslider0, "style", "knob");
        ui_interface -> addVerticalSlider("Volume", &fslider0,
           0.5f, 0.0f, 1.0f, 0.1f);
        ui_interface -> closeBox();
    virtual void compute (int count, FAUSTFLOAT ** input,
       FAUSTFLOAT** output) {
        float fSlow0 = (4.656613e-10f * float(fslider0));
        FAUSTFLOAT* output0 = output[0];
        for (int i=0; i < count; i++) {</pre>
            iRec0[0] = ((1103515245 * iRec0[1]) + 12345);
            output0[i] = (FAUSTFLOAT)(fSlow0 * iRec0[0]);
            // post processing
            iRec0[1] = iRec0[0];
        }
    }
};
int main(int argc, char* argv[])
{
mydsp DSP;
```

```
PrintUI ui;

// Activate the UI

// (here that only prints the control paths)
DSP.buildUserInterface(&ui);

// Allocate the audio driver to render

// 5 buffers of 512 frames
dummyaudio audio(5);
audio.init("Test", &DSP);

// Render buffers...
audio.start();

audio.stop();
}
```

You can possibly add the -i option to actually inline all #include "faust/xxx/yyy" headers (all files starting with "faust"). Then you will have to write a faust2xxx script that will chain the FAUST compilation step and the C++ compilation one. Look at scripts in the tools/faust2appls folder for real examples.

Developing the adapted C++ file may require "aggregating" the generated mydsp class (subclass of dsp base class defined in faust/dsp/dsp.h header) in your specific class, or "subclassing" and extend it. So you will have to write something like:

```
{
    // Do something specific
    fDSP.compute(count, inputs, outputs);
}

// Do something specific
};
or:
```

```
class my_class : public mydsp {
  private:
    // Do something specific
  public:
    my_class()
        // Do something specific
    virtual ~my_class()
        // Do something specific
    // Do something specific
    void my_compute(int count,
        FAUSTFLOAT ** inputs,
        FAUSTFLOAT** outputs,...)
    {
        // Do something specific
        compute(count, inputs, outputs);
    // Do something specific
};
```

This way your architecture file will be adapted to any "shape" of the generated code. That is, depending if you generate purely scalar, or vector code (using the -vec option), or any other option, the generated mydsp class will always be correctly inserted in the final C++ file. Look for instance at csound.cpp and unity.cpp architecture files in the architecture folder for real examples.

Chapter 7

OSC support

Most Faust architectures provide Open Sound Control (OSC) support ¹. This allows Faust applications to be remotely controlled from any OSC capable application, programming language, or hardware device. OSC support can be activated using the -osc option when building the application with the appropriate faust2xxx command. The following table (table 7.1) lists Faust's architectures which provide OSC support.

7.1 A simple example

To illustrate how OSC support works let's define a very simple noise generator with a level control: noise.dsp

```
process = library("music.lib").noise
    * hslider("level", 0, 0, 1, 0.01);
```

We are going to compile this example as a standalone Jack QT application with OSC support using the command:

```
faust2jaqt -osc noise.dsp
```

When we start the application from the command line:

```
./noise
```

we get various information on the standard output, including:

```
Faust OSC version 0.93 application 'noise' is running on UDP ports 5510, 5511, 5512
```

¹The implementation is based internally on the *oscpack* library by Ross Bencina

Audio system	Environment	OSC support
Linux		
Alsa	GTK, Qt, Console	yes
Jack	GTK, Qt, Console	yes
Netjack	GTK, Qt, Console	yes
PortAudio	GTK, Qt	yes
Mac OS X		
CoreAudio	Qt	yes
Jack	Qt, Console	yes
Netjack	Qt, Console	yes
PortAudio	Qt	yes
Windows		
Jack	Qt, Console	yes
PortAudio	Qt	yes

Table 7.1: FAUST architectures with OSC support.

As we can see the OSC module makes use of three different UDP ports:

- 5510 is the listening port number: control messages should be addressed to this port.
- 5511 is the output port number: control messages sent by the application and answers to query messages are sent to this port.
- 5512 is the error port number: used for asynchronous error notifications.

These OSC parameters can be changed from the command line using one of the following options:

- -port number set the port number used by the application to receive messages.
- -outport number set the port number used by the application to transmit messages.
- -errport number set the port number used by the application to transmit error messages.

- -desthost host set the destination host for the messages sent by the application.
- -xmit 0|1|2 turn transmission OFF, ALL, or ALIAS (default OFF). When transmission is OFF, input elements can be controlled using their addresses or aliases (if present). When transmission is ALL, input elements can be controlled using their addresses or aliases (if present), user's actions and output elements (bargraph) are transmitted as OSC messages as well as aliases (if present). When transmission is ALIAS, input elements can only be controlled using their aliases, user's actions and output elements (bargraph) are transmitted as aliases only.
- -xmitfilter path allows to filter output messages. Note that 'path' can be a regular expression (like "/freeverb/Reverbi/*").

For example:

```
./noise -xmit 1 -desthost 192.168.1.104 -outport 6000
```

will run noise with transmission mode ON, using 192.168.1.104 on port 6000 as destination.

7.2 Automatic port allocation

In order to address each application individually, only one application can be listening on a single port at one time. Therefore when the default incoming port 5510 is already opened by some other application, an application will automatically try increasing port numbers until it finds an available port. Let's say that we start two applications noise and mixer on the same machine, here is what we get:

```
$ ./noise &
...
Faust OSC version 0.93 application 'noise' is running
    on UDP ports 5510, 5511, 5512
$ ./mixer
...
Faust OSC version 0.93 application 'mixer' is running
    on UDP ports 5513, 5511, 5512
```

The mixer application fails to open the default incoming port 5510 because it is already opened by noise. Therefore it tries to find an available port starting from 5513 and open it. Please note that the two outcoming ports 5511 and 5512 are shared by all running applications.

7.3 Discovering OSC applications

oscsend hostname port address types values: send OpenSound Control message via UDP. types is a string, the letters indicates the type of the following values: i=integer, f=float, s=string,...

oscdump port : receive OpenSound Control messages via UDP and dump to standard output The commands oscsend Send OpenSound Control message via UDP. and oscdump-from the liblo package provide a convenient mean to experiment with OSC control. For the experiment let's use two additional terminals. The first one will be used to send OSC messages to the noise application using oscsend. The second terminal will be used to monitor the messages sent by the application using oscdump. We will indicate by T1\$ the command types on terminal T1 and by T2: the messages received on terminal T2. To monitor on terminal T2 the OSC messages received on UDP port 5511 we will use oscdump:

```
T2$ oscdump 5511
```

Once set we can use the hello message to scan UDP ports for FAUST applications. For example:

```
T1$ oscsend localhost 5510 "/*" s hello
```

gives us the root message address, the network and the UDP ports used by the noise application:

```
T2: /noise siii "192.168.1.102" 5510 5511 5512
```

7.4 Discovering the OSC interface of an application

Once we have an application we can discover its OSC interface (the set of OSC messages we can use to control it) by sending the get message to the root:

```
T1$ oscsend localhost 5510 /noise s get
```

As an answer of the osc messages understood by the application, a full description is available on terminal T2:

```
T2: /noise sF "xmit" #F
T2: /noise ss "desthost" "127.0.0.1"
T2: /noise si "outport" 5511
T2: /noise si "errport" 5512
T2: /noise/level fff 0.000000 0.000000 1.000000
```

The root of the osc interface is /noise. Transmission is OFF, xmit is set to false. The destination host for sending messages is "127.0.0.1", the output port is 5511 and the error port is 5512. The application has only one user interface element: /noise/level with current value 0.0, minimal value 0.0 and maximal value 1.0.

7.5 Widget's OSC address

Each widget of an application has a unique OSC address obtained by concatenating the labels of it's surrounding groups with its own label. Here is as an example mix4.dsp, a very simplified monophonic audio mixer with 4 inputs and one output. For each input we have a mute button and a level slider:

```
input(v) = vgroup("input %v", *(1-checkbox("mute")) :
    *(vslider("level", 0, 0, 1, 0.01)));
process = hgroup("mixer", par(i, 4, input(i)) :> _);
```

If we query this application:

```
T1$ oscsend localhost 5510 "/*" s get
```

We get a full description of its OSC interface on terminal T2:

```
T2: /mixer sF "xmit" #F

T2: /mixer ss "desthost" "127.0.0.1"

T2: /mixer si "outport" 5511

T2: /mixer si "errport" 5512

T2: /mixer/input_0/level fff 0.0000 0.0000 1.0000

T2: /mixer/input_0/mute fff 0.0000 0.0000 1.0000

T2: /mixer/input_1/level fff 0.0000 0.0000 1.0000

T2: /mixer/input_1/mute fff 0.0000 0.0000 1.0000

T2: /mixer/input_1/mute fff 0.0000 0.0000 1.0000

T2: /mixer/input_2/level fff 0.0000 0.0000 1.0000

T2: /mixer/input_2/mute fff 0.0000 0.0000 1.0000

T2: /mixer/input_3/level fff 0.0000 0.0000 1.0000

T2: /mixer/input_3/level fff 0.0000 0.0000 1.0000
```

As we can see each widget has a unique OSC address obtained by concatenating the top level group label "mixer", with the "input" group label and the widget label. Please note that in this operation whites spaces are replaced by underscores and metadata are removed.

All addresses must have a common root. This is the case in our example because there is a unique horizontal group "mixer" containing all widgets. If a common root is missing as in the following code:

```
input(v) = vgroup("input %v", *(1-checkbox("mute")) :
    *(vslider("level", 0, 0, 1, 0.01)));
process = par(i, 4, input(i)) :> _;
```

then a default vertical group is automatically create by the FAUST compiler using the name of the file mix4 as label:

There are potential conflicts between widget's labels and the OSC address space. An OSC symbolic name is an ASCII string consisting of a restricted set of printable characters. Therefore to ensure compatibility spaces are replaced by underscores and some other characters (asterisk, comma, forward, question mark, open bracket, close bracket, open curly brace, close curly brace) are replaced by hyphens.

```
T2: /mix4 sF "xmit" #F

T2: /mix4 ss "desthost" "127.0.0.1"

T2: /mix4 si "outport" 5511

T2: /mix4 si "errport" 5512

T2: /mix4/input_0/level fff 0.0000 0.0000 1.0000

T2: /mix4/input_0/mute fff 0.0000 0.0000 1.0000

T2: /mix4/input_1/level fff 0.0000 0.0000 1.0000

T2: /mix4/input_1/mute fff 0.0000 0.0000 1.0000

T2: /mix4/input_1/mute fff 0.0000 0.0000 1.0000

T2: /mix4/input_2/level fff 0.0000 0.0000 1.0000

T2: /mix4/input_2/mute fff 0.0000 0.0000 1.0000

T2: /mix4/input_3/level fff 0.0000 0.0000 1.0000

T2: /mix4/input_3/level fff 0.0000 0.0000 1.0000
```

7.6 Controlling the application via OSC

We can control any user interface element of the application by sending one of the previously discovered messages. For example to set the noise level of the application to 0.2 we send:

```
T1$ oscsend localhost 5510 /noise/level f 0.2
```

If we now query /noise/level we get, as expected, the value 0.2:

```
T1$ oscsend localhost 5510 /noise/level s get
T2: /noise/level fff 0.2000 0.0000 1.0000
```

7.7 Turning transmission ON

The xmit message at the root level is used to control the realtime transmission of OSC messages corresponding to user interface's actions. For examples:

```
T1$ oscsend localhost 5510 /noise si xmit 1
```

turns transmission in ALL mode. Now if we move the level slider we get a bunch of messages:

```
T2: /noise/level f 0.024000
T2: /noise/level f 0.032000
T2: /noise/level f 0.105000
T2: /noise/level f 0.250000
T2: /noise/level f 0.258000
T2: /noise/level f 0.185000
```

```
T2: /noise/level f 0.145000
T2: /noise/level f 0.121000
T2: /noise/level f 0.105000
T2: /noise/level f 0.008000
T2: /noise/level f 0.000000
```

This feature can be typically used for automation to record and replay actions on the user interface, or to remote control from one application to another. It can be turned OFF any time using:

```
T1$ oscsend localhost 5510 /noise si xmit 0
```

Use the ALIAS (xmit = 2) mode if you need restricted access to your program: when ALIAS is mode is used, only aliases of input elements (sliders, buttons...) can be used to control them, and output elements (bargraph) will only emit on their aliases.

7.8 Filtering OSC messages

When the transmission of OSC messages is ON, all the user interface elements are sent through the OSC connection.

```
T2: /harpe/level f 0.024000
T2: /harpe/hand f 0.1
T2: /harpe/level f 0.024000
T2: /harpe/hand f 0.25
T2: /harpe/level f 0.024000
T2: /harpe/hand f 0.44
T2: /noise/level f 0.145000
T2: /harpe/hand f 0.78
T2: /noise/level f 0.145000
T2: /harpe/hand f 0.99
```

We can choose to filter the unwanted parameters (or group of parameters). For example:

```
T1$ oscsend localhost 5510 /harpe si xmit 1 xmitfilter /harpe/level
```

As a result, we will receive:

```
T2: /harpe/hand f 0.1
T2: /harpe/hand f 0.25
T2: /harpe/hand f 0.44
T2: /harpe/hand f 0.78
```

To reset the filter, send:

```
T1$ oscsend localhost 5510 /harpe si xmit 1 xmitfilter
```

7.9 Using OSC aliases

Aliases are a convenient mechanism to control a FAUST application from a preexisting set of OSC messages.

Let's say we want to control our noise example with touchOSC on Android. The first step is to configure TouchOSC host to 192.168.1.102 (the host running our noise application) and outgoing port to 5510.

Then we can use oscdump 5510 (after quitting the noise application in order to free port 5510) to visualize the OSC messages sent by TouchOSC. Let's use for that the left slider of simple layout. Here is what we get:

```
T2: /1/fader1 f 0.000000
T2: /1/fader1 f 0.004975
T2: /1/fader1 f 0.004975
T2: /1/fader1 f 0.008125
T2: /1/fader1 f 0.017473
T2: /1/fader1 f 0.032499
T2: /1/fader1 f 0.051032
T2: ...
T2: /1/fader1 f 0.993289
T2: /1/fader1 f 1.000000
```

We can associate this OSC message to the noise level slider by inserting the metadata [osc:/1/fader1 o 1] into the slider's label:

```
process = library("music.lib").noise * hslider("level
    [osc:/1/fader1 0 1]",0,0,1,0.01);
```

Because here the range of /1/fader1 is 0 to 1 like the level slider we can remove the range mapping information and write simply:

```
process = library("music.lib").noise * hslider("level
       [osc:/1/fader1]", 0, 0, 1, 0.01);
```

TouchOSC can also send accelerometer data by enabling Settings/Options/Accelerometer. Using again oscdump 5510 we can visualize the messages send by TouchOSC:

```
T2: ...
T2: /accxyz fff -0.147842 0.019752 9.694721
T2: /accxyz fff -0.157419 0.016161 9.686341
```

Several osc aliases can be inserted into a single label allowing the same widget to be controlled by several OSC messages.

```
T2: /accxyz fff -0.167594 0.012570 9.683948
T2: ...
```

As we can see TouchOSC send the x, y and z accelerometers in a single message, as a triplet of values ranging approximatively from -9.81 to 9.81. In order to select the appropriate accelerometer we need to concatenate to /accxyz a suffix /0, /1 or /2. For example /accxyz/0 will correspond to x, /accxyz/1 to y, etc. We also need to define a mapping because the ranges are different:

```
process = library("music.lib").noise * hslider("level[
  osc:/accxyz/0 0 9.81]",0,0,1,0.01);
```

alias	description
[osc:/1/rotary1 0 1]	top left rotary knob
[osc:/1/rotary2 0 1]	middle left rotary knob
[osc:/1/rotary3 0 1]	bottom left rotary knob
[osc:/1/push1 0 1]	bottom left push button
[osc:/1/push2 0 1]	bottom center left push button
[osc:/1/toggle1 0 1]	top center left toggle button
[osc:/1/toggle2 0 1]	middle center left toggle button
[osc:/1/fader1 0 1]	center left vertical fader
[osc:/1/toggle3 0 1]	top center right toggle button
[osc:/1/toggle4 0 1]	middle center right toggle button
[osc:/1/fader2 0 1]	center right vertical toggle button
[osc:/1/rotary4 0 1]	top right rotary knob
[osc:/1/rotary5 0 1]	middle right rotary knob
[osc:/1/rotary6 0 1]	bottom right rotary knob
[osc:/1/push3 0 1]	bottom center right push button
[osc:/1/push4 0 1]	bottom right push button
[osc:/1/fader3 0 1]	bottom horizontal fader
[osc:/accxyz/0 -10 10]	x accelerometer
[osc:/accxyz/1 -10 10]	y accelerometer
[osc:/accxyz/2 -10 10]	z accelerometer

Table 7.2: Examples of OSC message aliases for TouchOSC (layout Mix2).

7.10 OSC cheat sheet

Default ports

```
default listening port
default transmission port
default error port
alternative listening ports
```

Command line options

```
-port n set the port number used by the application to receive messages set the port number used by the application to transmit messages -errport n set the port number used by the application to transmit error messages -desthost h set the destination host for the messages sent by the application -xmit 0|1|2 turn transmission OFF, ALL or ALIAS (default OFF) filter the FAUST paths at emission time
```

Discovery messages

```
oscsend host port "/*" s hello discover if any OSC application is listening on port p oscsend host port "/*" s get query OSC interface of application listening on port p
```

Control messages

```
oscsend host port "/*" si xmit 0|1|2 set transmission mode oscsend host port widget s get get widget's value oscsend host port widget f v set widget's value
```

Alias

```
"...[osc: address lo hi]..." alias with lo \rightarrow min, hi \rightarrow max mapping "...[osc: address]..." alias with min, max clipping
```

Chapter 8

HTTP support

Similarly to OSC, several Faust architectures also provide HTTP support. This allows Faust applications to be remotely controlled from any Web browser using specific URLs. Moreover OSC and HTTPD can be freely combined.

While OSC support is installed by default when FAUST is build, this is not the case for HTTP. That's because it depends on GNU *libmicrohttpd* library which is usually not installed by default on the system. An additional make httpd step is therefore required when compiling and installing FAUST:

```
make httpd
make
sudo make install
```

Note that make httpd will fail if *libmicrohttpd* is not available on the system.

The HTTP support can be activated using the -httpd option when building the audio application with the appropriate faust2xxx command. The following table (table 8.1) lists FAUST's architectures which provide HTTP support.

8.1 A simple example

To illustrate how HTTP support works let's reuse our previous mix4.dsp example, a very simplified monophonic audio mixer with 4 inputs and one output. For each input we have a mute button and a level slider:

Audio system	Environment	HTTP support
Linux		
Alsa	GTK, Qt, Console	yes
Jack	GTK, Qt, Console	yes
Netjack	GTK, Qt, Console	yes
PortAudio	GTK, Qt	yes
14 OG W		
Mac OS X		
CoreAudio	Qt	yes
Jack	Qt, Console	yes
Netjack	Qt, Console	yes
PortAudio	Qt	yes
Windows		
Jack	Qt, Console	yes
PortAudio	Qt	yes

Table 8.1: FAUST architectures with HTTP support.

We are going to compile this example as a standalone Jack QT application with HTTP support using the command:

```
faust2jaqt -httpd mix4.dsp
```

Th effect of the -httpd is to embed a small Web server into the application, which purpose is to serve an HTML page representing its user interface. This page makes use of JavaScript and SVG and is quite similar to the native QT interface.

When we start the application from the command line:

```
./mix4
```

we get various information on the standard output, including:

```
Faust httpd server version 0.72 is running on TCP port 5510
```

As we can see the embedded Web server is running by default on TCP port 5510. The entry point is http://localhost:5510. It can be open from any recent browser and it produces the page reproduced figure 8.1.

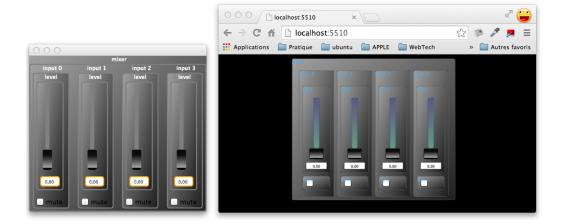


Figure 8.1: User interface of mix4.dsp in a Web browser

8.2 JSON description of the user interface

The communication between the application and the Web browser is based on several underlying URLs. The first one is http://localhost:5510/JSON that return a json description of the user interface of the application. This json description is used internally by the JavaScript code to build the graphical user interface. Here is (part of) the json returned by mix4:

```
{
  "name": "mix4",
  "address": "YannAir.local",
  "port": "5511",
  "ui": [
    {
      "type": "hgroup",
      "label": "mixer",
      "items": [
          "type": "vgroup",
          "label": "input_0",
          "items": [
               "type": "vslider",
              "label": "level",
              "address": "/mixer/input_0/level",
              "init": "0", "min": "0", "max": "1",
               "step": "0.01"
```

```
{
    "type": "checkbox",
    "label": "mute",
    "address": "/mixer/input_0/mute",
    "init": "0", "min": "0", "max": "0",
    "step": "0"
    }
    ...
}
```

8.3 Quering the state of the application

Each widget has a unique "address" field that can be used to query its value. In our example here the level of the input o has the address /mixer/input_0/level. The address can be used to forge an URL to get the value of the widget: http://localhost:5510/mixer/input_0/level, resulting in:

```
/mixer/input_0/level 0.00000
```

Multiple widgets can be query at once by using an address higher in the hierarchy. For example to get the values of the level and the mute state of input o we use http://localhost:5510/mixer/input_0, resulting in:

```
/mixer/input_0/level 0.00000
/mixer/input_0/mute 0.00000
```

To get the all the values at once we simply use http://localhost:5510/mixer, resulting in:

```
/mixer/input_0/level 0.00000
/mixer/input_0/mute 0.00000
/mixer/input_1/level 0.00000
/mixer/input_1/mute 0.00000
/mixer/input_2/level 0.00000
/mixer/input_2/mute 0.00000
/mixer/input_3/level 0.00000
/mixer/input_3/mute 0.00000
```

8.4 Changing the value of a widget

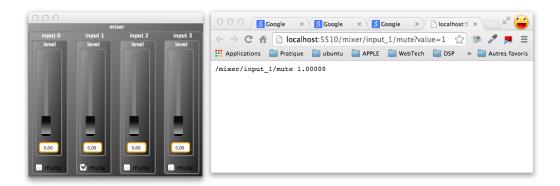


Figure 8.2: Muting input 1 by forging the appropriate URL

Let's say that we want to mute input I of our mixer. We can use for that purpose the URL http://localhost:5510/mixer/input_1/mute?value=1 obtained by concatenating?value=1 at the end of the widget URL.

All widgets can be controlled similarly. For example http://localhost:5510/mixer/input_3/level?value=0.7 will sets the input 3 level to 0.7.

8.5 Proxy control access to the Web server

A control application may want to access and control the running DSP using its Web server, but without using the delivered HTML page in a browser. Since the complete json can be retrieved, control applications can purely be developed in C/C++, then build a *proxy* version of the use interface, and set and get parameters using HTTP requests.

This mode can be started dynamically using the *-server URL* parameter. Assuming an application with HTTP support is running remotely on the given URL, the control application will fetch its json description, use it to dynamically build the user interface, and allow to access the remote parameters.

8.6 HTTP cheat sheet

Here is a summary of the various URLs used to interact with the application's Web server.

Default ports

```
default TCP port used by the application's Web server alternative TCP ports
```

Command line options

```
-port \,n\, set the TCP port number used by the application's Web server -server \,URL\, start a proxy control application accessing the remote application running on the give
```

URLs

```
http://host:port the base URL to be used in proxy control access mode
http://host:port/JSON get a json description of the user interface
http://host:port/address
http://host:port/address?value=v set the value of a widget to v

set the value of a widget to v
```

JSON

Top level

The json describes the name, host and port of the application and a hierarchy of user interface items:

```
{
    "name": <name>,
    "address": <host>,
    "port": <port>,
    "ui": [ <item> ]
}
```

An <item> is either a group (of items) or a widget.

Groups

A group is essentially a list of items with a specific layout:

```
{
    "type": <type>,
    "label": <label>,
    "items": [ <item>, <item>,...]
}
```

The <type> defines the layout. It can be either "vgroup", "hgroup" or "tgroup"

Widgets

```
{
    "type": <type>,
    "label": <label>,
    "address": <address>,
    "meta": [ { "key": "value"},...],
    "init": <num>,
    "min": <num>,
    "max": <num>,
    "step": <num>)
},
```

Widgets are the basic items of the user interface. They can be of different <type>: "button", "checkbox", "nentry", "vslider", "hslider", "vbargraph" or "hbargraph".

Chapter 9

MIDI support

Similarly to OSC, several FAUST architectures also provide MIDI support. This allows FAUST applications to be controlled from any MIDI device (or to control MIDI devices). MIDI is also the preferable way to control Polyphonic instruments.

9.1 MIDI messages description in the dsp source code

MIDI control messages are described as metadata in UI elements. They are decoded by a special architecture *MidiUI* class that will parse incoming MIDI messages and update the appropriate control parameters, or send MIDI messages when the UI elements (sliders, buttons...) are moved.

9.2 Description of the possible standard MIDI messages

Below, when a 7-bit MIDI parameter is used to drive a button or checkbox, its maximum value 127 maps to 1 ("on") while its minimum value 0 maps to 0 ("off").

A special [midi:xxx yyy...] metadata needs to be added in the UI element description. The more usual MIDI messages can be used as described here:

- [midi:ctrl num] in a slider or bargraph will map the UI element value to (0, 127) range. When used with a button or checkbox, I will be mapped to 127, 0 will be mapped to 0,
- [midi:keyon pitch] in a slider or bargraph will register the UI element's state-variable to be driven by MIDI note-on velocity (an integer between 0 and 127) of the specified key between 0 and 127. When used with a button or checkbox, I will be mapped to 127, 0 will be mapped to 0,

- [midi:keyoff pitch] in a slider or bargraph will register the UI element's state-variable to be driven by MIDI note-off velocity (an integer between 0 and 127) of the specified key between 0 and 127. When used with a button or checkbox, I will be mapped to 127, 0 will be mapped to 0,
- [midi:key pitch] in a slider or bargraph will register the UI element's state-variable to be driven by MIDI note-on velocity (an integer between 0 and 127) of the specified key between 0 and 127. When used with a button or checkbox, I will be mapped to 127, 0 will be mapped to 0. Note-on and note-off events will be handled,
- [midi:keypress pitch] in a slider or bargraph will register the UI element's statevariable to be driven by the MIDI key-pressure (an integer between o and 127) from MIDI key,
- [midi:pgm num] in a slider or bargraph will map the UI element value to the progchange value, so *progchange* message with the same *num* value will be sent. When used with a button or checkbox, I will send the *progchange* message with *num* value, o will send nothing,
- [midi: chanpress num] in a slider or bargraph will map the UI element value to the chanpress value, so *chanpress* message with the same *num* value will be sent. When used with a button or checkbox, I will send the *chanpress* message with *num* value, o will send nothing,
- [midi:pitchwheel] in a slider or bargraph will map the UI element value to (0,16383) range. When used with a button or checkbox, 1 will be mapped to 16383, 0 will be mapped to 0.

9.3 A simple example

An example with a *volume* slider controlled with MIDI ctrlchange 7 messages :

A complete testing example named *midi_tester.dsp* is available in the FAUST distribution *examples* folder.



Figure 9.1: MIDI messages testing example

The MIDI support can be activated using the -midi option when building the audio application with the appropriate faust2xxx command. The following table (table 9.1) lists FAUST's architectures which provide MIDI support.

Audio system	Environment	HTTP support
<i>Linux</i> Alsa Jack	Qt Qt	yes yes
<i>Mac OS X</i> CoreAudio Jack	Qt Qt	yes yes

Table 9.1: FAUST architectures with HTTP support.

9.4 MIDI synchronization

MIDI clock based synchronization can be used to slave a given Faust program. The following three messages need to be used:

- [midi:start] in a button or checkbox will trigger a value of 1 when a *start* MIDI message is received
- [midi:stop] in a button or checkbox will trigger a value of o when a *stop* MIDI message is received
- [midi:clock] in a button or checkbox will deliver a sequence of successive 1 and 0 values each time a *clock* MIDI message is received, seen by FAUST code as a square command signal, to be used to compute higher level information.

A typical Faust program will then use the MIDI clock stream to possibly compute the BPM information, or for any synchronization need it may have. Here is a simple example of a sinus generated which a frequency controlled by the MIDI clock stream, and starting/stopping when receiving the MIDI start/stop messages:

```
import("music.lib");

// square signal (1/0), changing state at each received clock
clocker = checkbox("MIDI clock[midi:clock]");

// ON/OFF button controlled with MIDI start/stop messages
play = checkbox("ON/OFF [midi:start] [midi:stop]");

// detect front
front(x) = (x-x') != 0.0;

// count number of peaks during one second
freq(x) = (x-x@SR) : + ~ _;

process = osc(8*freq(front(clocker))) * play;
```

Chapter 10

Polyphonic support

Directly programing polyphonic instruments in FAUST is perfectly possible. It is also needed if very complex signal interaction between the different voices have to be described¹.

But since all voices would always be computed, this approach could be too CPU costly for simpler or more limited needs. In this case describing a single voice in a FAUST DSP program and externally combining several of them with a special *polyphonic instrument aware* architecture file is a better solution. Moreover, this special architecture file takes care of dynamic voice allocations and control MIDI messages decoding and mapping.

10.1 Polyphonic ready DSP code

By convention FAUST architecture files with polyphonic capabilities expect to find control parameters named *freq*, *gain* and *gate*. The metadata declare nvoices "8"; kind of line with a desired value of voices can be added in the source code.

In the case of MIDI control, the *freq* parameter (which should be a frequency) will be automatically computed from MIDI note numbers, *gain* (which should be a value between o and 1) from velocity and *gate* from *keyon/keyoff* events. Thus, gate can be used as a trigger signal for any envelope generator, etc.

10.2 Using the mydsp_poly class

The single voice has to be described by a FAUST DSP program, the mydsp_poly class is then used to combine several voices and create a polyphonic ready DSP:

¹Like sympathetic strings resonance in a physical model of a piano for instance.

- the *faust/dsp/poly-dsp.h* file contains the definition of the mydsp_poly class used to wrap the DSP voice into the polyphonic architecture. This class maintains an array of dsp type of objects, manage dynamic voice allocations, control MIDI messages decoding and mapping, mixing of all running voices, and stopping a voice when its output level decreases below a given threshold.
- as a sub-class of DSP, the mydsp_poly class redefines the buildUserInterface method. By convention all allocated voices are grouped in a global *Polyphonic* tabgroup. The first tab contains a *Voices* group, a master like component used to change parameters on all voices at the same time, with a *Panic* button to be used to stop running voices², followed by one tab for each voice. Graphical User Interface components will then reflect the multi-voices structure of the new polyphonic DSP (Figure 10.1).

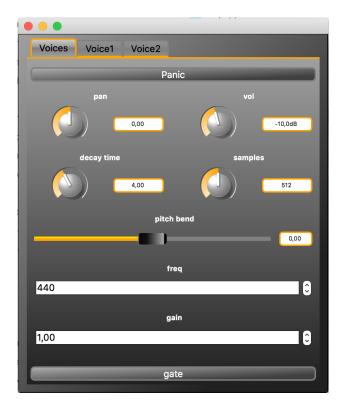


Figure 10.1: Extended multi-voices GUI interface

The resulting polyphonic DSP object can be used as usual, connected with the needed audio driver, and possibly other UI control objects like OSCUI, httpdUI, etc. Having

²An internal control grouping mechanism has been defined to automatically dispatch a user interface action done on the master component on all linked voices, except for the *freq*, *gain* and *gate* controls.

this new UI hierarchical view allows complete OSC control of each single voice and their control parameters, but also all voices using the master component.

The following OSC messages reflect the same DSP code either compiled normally, or in polyphonic mode (only part of the OSC hierarchies are displayed here):

```
// Mono mode

/0x00/0x00/vol f -10.0
/0x00/0x00/pan f 0.0

// Polyphonic mode

/Polyphonic/Voices/0x00/0x00/pan f 0.0
/Polyphonic/Voices/0x00/0x00/vol f -10.0
...
/Polyphonic/Voice1/0x00/0x00/vol f -10.0
/Polyphonic/Voice1/0x00/0x00/pan f 0.0
...
/Polyphonic/Voice2/0x00/0x00/vol f -10.0
/Polyphonic/Voice2/0x00/0x00/pan f 0.0
...
```

The polyphonic instrument allocation takes the DSP to be used for one voice³, the desired number of voices, the *dynamic voice allocation* state⁴, and the *group* state which controls if separated voices are displayed or not (Figure 10.1):

```
DSP = new mydsp_poly(dsp, 2, true, true);
```

With the following code, note that a polyphonic instrument may be used outside of a MIDI control context, so that all voices will be always running and possibly controlled with OSC messages for instance:

```
DSP = new mydsp_poly(dsp, 8, false, true);
```

10.3 Controlling the polyphonic instrument

The mydsp_poly class is also ready for MIDI control and can react to *keyon/keyoff* and *pitchwheel* messages. Other MIDI control parameters can directly be added in the DSP source code.

³The DSP object will be automatically cloned in the mydsp_poly class to create all needed voices.

⁴Voices may be always running, or dynamically started/stopped in case of MIDI control.

10.4 Deploying the polyphonic instrument

Several architecture files and associated scripts have been updated to handle polyphonic instruments:

As an example on OSX, the script faust2caqt foo.dsp can be used to create a polyphonic CoreAudio/QT application. The desired number of voices is either declared in a nvoices metadata or changed with the -nvoices num additional parameter. MIDI control is activated using the -midi parameter.

The number of allocated voices can possibly be changed at runtime using the -nvoices parameter to change the default value (so using ./foo -nvoices 16 for instance).

Several other scripts have been adapted using the same conventions.

10.5 Polyphonic instrument with a global output effect

Polyphonic instruments may be used with an output effect. Putting that effect in the main FAUST code is not a good idea since it would be instantiated for each voice which would be very inefficient. This is a typical use case for the dsp_sequencer class previously presented with the polyphonic DSP connected in sequence with a unique global effect (Figure 10.2).

faustcaqt inst.dsp -effect effect.dsp with inst.dsp and effect.dsp in the same folder, and the number of outputs of the instrument matching the number of inputs of the effect, has to be used. A dsp_sequencer object will be created to combine the polyphonic instrument in sequence with the single output effect.

Polyphonic ready *faust2xx* scripts will then compile the polyphonic instrument and the effect, combine them in sequence, and create a ready to use DSP.

10.5.1 Integrated global output effect

Starting with the 2.5.17 version, a new convention has been defined to directly integrate a global output effect inside the DSP source code itself. The effect has simply to be declared in a effect = effect_code; line in the source. Here is a more complete source code example:

```
import("stdfaust.lib");
process = pm.clarinet_ui_MIDI <: _,_;
effect = dm.freeverb_demo;</pre>
```

⁵⁻nvoices parameter takes precedence over the metadata value.

10.5. POLYPHONIC INSTRUMENT WITH A GLOBAL OUTPUT EFFECT 101



Figure 10.2: Polyphonic instrument with output effect GUI interface: left tab window shows the polyphonic instrument with its *Voices* group only, right tab window shows the output effect.

The architecture script then separates the instrument description itself (the **process** = ... definition) from the effect definition (the effect = ... definition), possibly adapts the instrument number of outputs to the effect number of inputs, compiles each part separately, and combines them with the dsp_sequencer object.

A new auto parameter to be used in *faust2xx* script has been defined, as in the faustcaqt inst.dsp -effect auto line for example.

10.5.2 Integrated global output effect and libfaust

For developers using the libfaust library, an helper file named faust/dsp/poly-dsp-tools.h is available. It defines an API to automatically create a polyphonic instrument with an output effect, starting from a DSP source file using the effect effect = ... convention. The function createPolyDSPFactoryFromString or createPolyDSPFactoryFromFile must be used to create the polyphonic DSP factory. Next, the createPolyDSPInstance function creates the polyphonic object (a subclass of dsp_poly type) to be used like a regular dsp type object.

After the DSP factory has been compiled, your application or plugin may want to save/restore it in order to save FAUST to LLVM IR compilation or even JIT compi-

lation time at next use. To get the internal factory compiled code, several functions are available:

- writePolyDSPFactoryToIRFile allows to save the polyphonic factory LLVM IR (in textual format) in a file,
- writePolyDSPFactoryToBitcodeFile allows to save the polyphonic factory LLVM IR (in binary format) in a file,
- writePolyDSPFactoryToMachineFile allows to save the polyphonic factory executable machine code in a file.

To re-create a DSP factory from a previously saved code, several functions are available:

- readPolyDSPFactoryFromIRFile allows to create a polyphonic DSP factory from a file containing the LLVM IR (in textual format),
- readPolyDSPFactoryFromBitcodeFile allows to create a polyphonic factory from a file containing the LLVM IR (in binary format),
- readPolyDSPFactoryFromMachineFile allows to create a polyphonic DSP factory from a file containing the executable machine code.

Chapter 11

Controlling the code generation

Several options of the FAUST compiler allow to control the generated C++ code. By default the computations are done sample by sample in a single loop. But the compiler can also generate *vector* and *parallel* code.

II.I Vector code generation

Modern C++ compilers are able to do autovectorization, that is to use SIMD instructions to speedup the code. These instructions can typically operate in parallel on short vectors of 4 simple precision floating point numbers thus leading to a theoretical speedup of $\times 4$. Autovectorization of C/C++ programs is a difficult task. Current compilers are very sensitive to the way the code is arranged. In particular too complex loops can prevent autovectorization. The goal of the vector code generation is to rearrange the C++ code in a way that facilitates the autovectorization job of the C++ compiler. Instead of generating a single sample computation loop, it splits the computation into several simpler loops that communicates by vectors.

The vector code generation is activated by passing the --vectorize (or -vec) option to the FAUST compiler. Two additional options are available: --vec-size <n> controls the size of the vector (by default 32 samples) and --loop-variant 0/1 gives some additional control on the loops.

To illustrate the difference between scalar code and vector code, let's take the computation of the RMS (Root Mean Square) value of a signal. Here is the FAUST code that computes the Root Mean Square of a sliding window of 1000 samples:

```
// Root Mean Square of n consecutive samples
RMS(n) = square : mean(n) : sqrt;
```

The compute() method generated in scalar mode is the following:

```
virtual void compute (int count,
                       float ** input,
                       float** output)
  float* input0 = input[0];
  float* output0 = output[0];
  for (int i=0; i < count; i++) {</pre>
    float fTemp0 = input0[i];
    int iTemp1 = int(1048576*fTemp0*fTemp0);
    iVec0[IOTA\&1023] = iTemp1;
    iRec0[0] = ((iVec0[IOTA&1023] + iRec0[1])
                     - iVec0[(IOTA-1000)&1023]);
    output0[i] = sqrtf(9.536744e-10f *
                        float(iRec0[0]));
    // post processing
    iRec0[1] = iRec0[0];
    IOTA = IOTA+1;
  }
}
```

The -vec option leads to the following reorganization of the code:

```
float** output)
{
  int iRec0_{tmp}[32+4];
  int* iRec0 = &iRec0_tmp[4];
  for (int index=0; index<fullcount; index+=32)</pre>
    int count = min (32, fullcount-index);
    float* input0 = &input[0][index];
    float* output0 = &output[0][index];
    for (int i=0; i<4; i++)
      iRecO_tmp[i]=iRecO_perm[i];
    // SECTION : 1
    for (int i=0; i < count; i++) {
      iYec0[(iYec0_idx+i)&2047] =
               int(1048576*input0[i]*input0[i]);
    }
    // SECTION : 2
    for (int i=0; i < count; i++) {</pre>
      iRec0[i] = ((iYec0[i] + iRec0[i-1]) -
                iYec0[(iYec0_idx+i-1000)&2047]);
    }
    // SECTION : 3
    for (int i=0; i < count; i++) {</pre>
      output0[i] = sqrtf((9.536744e-10f *
                  float(iRec0[i])));
    }
    // SECTION : 4
    iYec0_idx = (iYec0_idx+count)&2047;
    for (int i=0; i<4; i++)
      iRec0_perm[i]=iRec0_tmp[count+i];
  }
```

While the second version of the code is more complex, it turns out to be much easier to vectorize efficiently by the C++ compiler. Using Intel icc II.0, with the exact same compilation options: -03 -xHost -ftz -fno-alias -fp-model fast=2, the scalar version leads to a throughput performance of I29.I44 MB/s, while the vector version achieves 359.548 MB/s, a speedup of x2.8!

The vector code generation is built on top of the scalar code generation (see figure 11.1). Every time an expression needs to be compiled, the compiler checks if it requires a separate loop or not. It applies some simple rules for that. Expressions that are shared (and are complex enough) are good candidates to be compiled in a separate loop, as well as recursive expressions and expressions used in delay lines.

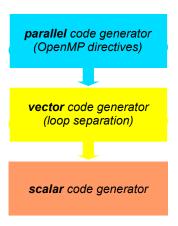


Figure 11.1: FAUST's stack of code generators

The result is a directed graph in which each node is a computation loop (see Figure II.2). This graph is stored in the klass object and a topological sort is applied to it before printing the code.

11.2 Parallel code generation

The parallel code generation is activated by passing either the --openMP (or -omp) option or the --scheduler (or -sch) option. It implies the -vec options as the parallel code generation is built on top of the vector code generation.

II.2.1 The OpenMP code generator

The --openMP (or -omp) option given to the FAUST compiler will insert appropriate OpenMP directives in the C++ code. OpenMP (http://wwww.openmp.org) is a well established API that is used to explicitly define direct multi-threaded, shared memory parallelism. It is based on a fork-join model of parallelism (see figure II.3). Parallel regions are delimited by #pragma omp parallel constructs. At the entrance of a parallel region a team of parallel threads is activated. The code within a parallel region is executed by each thread of the parallel team until the end of the region.

```
#pragma omp parallel
{
```

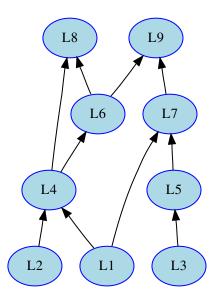


Figure 11.2: The result of the -vec option is a directed acyclic graph (DAG) of small computation loops

```
// the code here is executed simultaneously by
// every thread of the parallel team
...
}
```

In order not to have every thread doing redundantly the exact same work, OpemMP provides specific *work-sharing* directives. For example #pragma omp sections allows to break the work into separate, discrete sections, each section being executed by one thread:

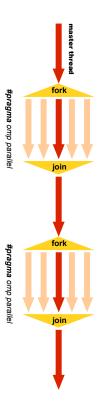


Figure 11.3: OpenMP is based on a fork-join model

```
// job 2
}
...
}
...
}
```

11.2.2 Adding OpenMP directives

As said before the parallel code generation is built on top of the vector code generation. The graph of loops produced by the vector code generator is topologically sorted in order to detect the loops that can be computed in parallel. The first set S_0 (loops L1, L2 and L3 in the DAG of Figure II.2) contains the loops that don't depend on any other loops, the set S_1 contains the loops that only depend on loops of S_0 , (that is loops L4 and L5), etc..

As all the loops of a given set S_n can be computed in parallel, the compiler will generate

a sections construct with a section for each loop.

```
#pragma omp sections
{
    #pragma omp section
    for (...) {
        // Loop 1
    }
    #pragma omp section
    for (...) {
        // Loop 2
    }
    ...
}
```

If a given set contains only one loop, then the compiler checks to see if the loop can be parallelized (no recursive dependencies) or not. If it can be parallelized, it generates:

```
#pragma omp for
for (...) {
   // Loop code
}
```

otherwise it generates a single construct so that only one thread will execute the loop:

```
#pragma omp single
for (...) {
   // Loop code
}
```

11.2.3 Example of parallel OpenMP code

To illustrate how FAUST uses the OpenMP directives, here is a very simple example, two 1-pole filters in parallel connected to an adder (see figure 11.4 the corresponding block-diagram):

```
filter(c) = *(1-c) : + ~ *(c);
process = filter(0.9), filter(0.9) : +;
```

The corresponding compute() method obtained using the -omp option is the following:

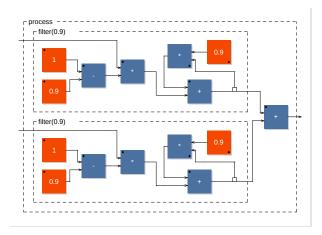


Figure 11.4: two filters in parallel connected to an adder

```
float
        fRec0_tmp[32+4];
        fRec1_tmp[32+4];
float
float* fRec0 = &fRec0_tmp[4];
        fRec1 = &fRec1_tmp[4];
#pragma omp parallel firstprivate(fRec0,fRec1)
  for (int index = 0; index < fullcount;</pre>
                               index += 32)
  {
    int count = min (32, fullcount-index);
    float* input0 = &input[0][index];
    float* input1 = &input[1][index];
    float* output0 = &output[0][index];
    #pragma omp single
      for (int i=0; i<4; i++)
        fRecO_tmp[i]=fRecO_perm[i];
      for (int i=0; i<4; i++)
        fRec1_tmp[i]=fRec1_perm[i];
    }
    // SECTION : 1
    #pragma omp sections
      #pragma omp section
      for (int i=0; i < count; i++) {</pre>
        fRec0[i] = ((0.1f * input1[i])
```

```
+ (0.9f * fRec0[i-1]));
      #pragma omp section
      for (int i=0; i < count; i++) {
        fRec1[i] = ((0.1f * input0[i])
                  + (0.9f * fRec1[i-1]));
      }
    }
    // SECTION : 2
    #pragma omp for
    for (int i=0; i < count; i++) {
      output0[i] = (fRec1[i] + fRec0[i]);
    // SECTION : 3
    #pragma omp single
      for (int i=0; i<4; i++)
        fRecO_perm[i]=fRecO_tmp[count+i];
      for (int i=0; i<4; i++)
        fRec1_perm[i]=fRec1_tmp[count+i];
    }
  }
}
```

This code requires some comments:

- I. The parallel construct #pragma omp parallel is the fundamental construct that starts parallel execution. The number of parallel threads is generally the number of CPU cores but it can be controlled in several ways.
- 2. Variables external to the parallel region are shared by default. The pragma firstprivate (fRec0,fRec1) indicates that each thread should have its private copy of fReco and fReci. The reason is that accessing shared variables requires an indirection and is quite inefficient compared to private copies.
- 3. The top level loop for (int index = 0;...)... is executed by all threads simultaneously. The subsequent work-sharing directives inside the loop will indicate how the work must be shared between the threads.
- 4. Please note that an implied barrier exists at the end of each work-sharing region. All threads must have executed the barrier before any of them can continue.

- 5. The work-sharing directive #pragma omp single indicates that this first section will be executed by only one thread (any of them).
- 6. The work-sharing directive #pragma omp sections indicates that each corresponding #pragma omp section, here our two filters, will be executed in parallel.
- 7. The loop construct #pragma omp for specifies that the iterations of the associated loop will be executed in parallel. The iterations of the loop are distributed across the parallel threads. For example, if we have two threads, the first one can compute indices between 0 and count/2 and the other one between count/2 and count.
- 8. Finally #pragma omp single in section 3 indicates that this last section will be executed by only one thread (any of them).

II.2.4 The scheduler code generator

With the --scheduler (or -sch) option given to the Faust compiler, the computation graph is cut into separated computation loops (called "tasks"), and a "Work Stealing Scheduler" is used to activate and execute them following their dependencies. A pool of worked threads is created and each thread uses it's own local WSQ (Work Stealing Queue) of tasks. A WSQ is a special queue with a Push operation, a "private" LIFO Pop operation and a "public" FIFO Pop operation.

Starting from a ready task, each thread follows the dependencies, possibly pushing ready sub-tasks into it's own local WSQ. When no more tasks can be activated on a given computation path, the thread pops a task from it's local WSQ. If the WSQ is empty, then the thread is allowed to "steal" tasks from other threads WSQ.

The local LIFO Pop operation allows better cache locality and the FIFO steal Pop "larger chuck" of work to be done. The reason for this is that many work stealing workloads are divide-and-conquer in nature, stealing one of the oldest task implicitly also steals a (potentially) large subtree of computations that will unfold once that piece of work is stolen and run.

Compared to the OpenMP model (-omp) the new model is worse for simple FAUST programs and usually starts to behave comparable or sometimes better for "complex enough" FAUST programs. In any case, since OpenMP does not behave so well with GCC compilers (only quite recent versions like GCC 4.4 start to show some improvements), and is unusable on OSX in real-time contexts, this new scheduler option has it's own value. We plan to improve it adding a "pipelining" idea in the future.

II.2.5 Example of parallel scheduler code

To illustrate how FAUST generates the scheduler code, here is a very simple example, two 1-pole filters in parallel connected to an adder (see figure 11.4 the corresponding block-diagram):

```
filter(c) = *(1-c) : + ~ *(c);
process = filter(0.9), filter(0.9) : +;
```

When -sch option is used, the content of the additional *architecture/scheduler.h* file is inserted in the generated code. It contains code to deal with WSQ and thread management. The compute() and computeThread() methods are the following:

```
virtual void compute (int fullcount,
                       float ** input,
                       float ** output)
{
    GetRealTime();
    this->input = input;
    this->output = output;
    StartMeasure();
    for (fIndex = 0; fIndex < fullcount; fIndex += 32) {</pre>
        fFullCount = min (32, fullcount-fIndex);
        TaskQueue::Init();
        // Initialize end task
        fGraph.InitTask(1,1);
        // Only initialize tasks with inputs
        fGraph.InitTask(4,2);
        fIsFinished = false;
        fThreadPool.SignalAll(fDynamicNumThreads - 1);
        computeThread(0);
        while (!fThreadPool.IsFinished()) {}
    StopMeasure(fStaticNumThreads,
        fDynamicNumThreads);
void computeThread (int cur_thread) {
    float* fRec0 = &fRec0_tmp[4];
    float* fRec1 = &fRec1_tmp[4];
    // Init graph state
        TaskQueue taskqueue;
        int tasknum = -1;
        int count = fFullCount;
```

```
// Init input and output
FAUSTFLOAT* input0 = &input[0][fIndex];
FAUSTFLOAT* input1 = &input[1][fIndex];
FAUSTFLOAT* output0 = &output[0][fIndex];
int task_list_size = 2;
int task_list[2] = {2,3};
taskqueue.InitTaskList(task_list_size, task_list,
    fDynamicNumThreads, cur_thread, tasknum);
while (!fIsFinished) {
    switch (tasknum) {
        case WORK_STEALING_INDEX: {
            tasknum = TaskQueue::GetNextTask(
               cur_thread);
            break;
        }
        case LAST_TASK_INDEX: {
            fIsFinished = true;
            break;
        }
        // SECTION : 1
        case 2: {
            // LOOP 0x101111680
            // pre processing
            for (int i=0; i<4; i++) fRec0_tmp[i]=
               fRecO_perm[i];
            // exec code
            for (int i=0; i < count; i++) {</pre>
                fRec0[i] = ((1.000000e-01f * (
                   float)input1[i]) + (0.9f *
                   fRec0[i-1]));
            }
            // post processing
            for (int i=0; i<4; i++) fRec0_perm[i
               ] = fRec0_tmp[count+i];
            fGraph.ActivateOneOutputTask(
               taskqueue, 4, tasknum);
            break;
        }
        case 3: {
            // LOOP 0x1011125e0
            // pre processing
            for (int i=0; i<4; i++) fRec1_tmp[i]=
```

```
fRec1_perm[i];
                 // exec code
                 for (int i=0; i < count; i++) {</pre>
                      fRec1[i] = ((1.000000e-01f * (
                         float)input0[i]) + (0.9f *
                         fRec1[i-1]));
                 }
                 // post processing
                 for (int i=0; i<4; i++) fRec1_perm[i</pre>
                     ]=fRec1_tmp[count+i];
                 fGraph.ActivateOneOutputTask(
                     taskqueue, 4, tasknum);
                 break;
             }
             case 4: {
                 // LOOP 0x101111580
                 // exec code
                 for (int i=0; i < count; i++) {</pre>
                      output0[i] = (FAUSTFLOAT)(fRec1[i
                         ] + fRec0[i]);
                 }
                 tasknum = LAST_TASK_INDEX;
                  break;
             }
        }
   }
}
```

Chapter 12

Mathematical Documentation

The FAUST compiler provides a mechanism to produce a self-describing documentation of the mathematical semantic of a FAUST program, essentially as a pdf file. The corresponding options are -mdoc (short) or --mathdoc (long).

12.1 Goals of the mathdoc

There are three main goals, or uses, of this mathematical documentation:

- 1. to preserve signal processors, independently from any computer language but only under a mathematical form;
- 2. to bring some help for debugging tasks, by showing the formulas as they are really computed after the compilation stage;
- 3. to give a new teaching support, as a bridge between code and formulas for signal processing.

12.2 Installation requirements

- faust, of course!
- svg2pdf (from the Cairo 2D graphics library), to convert block-diagrams, as LETEX doesn't eat SvG directly yet...
- breqn, a LTEX package to handle automatic breaking of long equations,
- pdflatex, to compile the LATEX output file.

12.3 Generating the mathdoc

The easiest way to generate the complete mathematical documentation is to call the faust2mathdoc script on a FAUST file, as the -mdoc option leave the documentation production unfinished. For example:

```
faust2mathdoc noise.dsp
```

12.3.1 Invoking the -mdoc option

Calling directly faust -mdoc does only the first part of the work, generating:

- a top-level directory, suffixed with "-mdoc",
- 5 subdirectories (cpp/, pdf/, src/, svg/, tex/),
- a LaTEX file containing the formulas,
- Svg files for block-diagrams.

At this stage:

- · cpp/ remains empty,
- pdf/remains empty,
- src/contains all FAUST sources used (even libraries),
- svg/ contains SvG block-diagram files,
- tex/ contains the generated LATEX file.

12.3.2 Invoking faust2mathdoc

The faust2mathdoc script calls faust --mathdoc first, then it finishes the work:

- moving the output C++ file into cpp/,
- converting all SvG files into pdf files (you must have svg2pdf installed, from the Cairo 2D graphics library),
- launching pdflatex on the ETEX file (you must have both pdflatex and the breqn package installed),
- moving the resulting pdf file into pdf/.

12.3.3 Online examples

To get an idea of the results of this mathematical documentation, which captures the mathematical semantic of FAUST programs, you can look at two pdf files online:

- http://faust.grame.fr/pdf/karplus.pdf (automatic documentation),
- http://faust.grame.fr/pdf/noise.pdf (manual documentation).

You can also generate all *mdoc* pdfs at once, simply invoking the make mathdoc command inside the examples/ directory:

- for each %. dsp file, a complete %-mdoc directory will be generated,
- a single allmathpdfs/directory will gather all the generated pdf files.

12.4 Automatic documentation

By default, when no <mdoc> tag can be found in the input FAUST file, the -mdoc option automatically generates a Lagrangian with four sections:

- I. "Equations of process", gathering all formulas needed for process,
- 2. "Block-diagram schema of process", showing the top-level block-diagram of process
- 3. "Notice of this documentation", summing up generation and conventions information,
- 4. "Complete listing of the input code", listing all needed input files (including libraries).

12.5 Manual documentation

You can specify yourself the documentation instead of using the automatic mode, with five xml-like tags. That permits you to modify the presentation and to add your own comments, not only on process, but also about any expression you'd like to. Note that as soon as you declare an <mdoc> tag inside your FAUST file, the default structure of the automatic mode is ignored, and all the ETEX stuff becomes up to you!

12.5.1 Six tags

Here are the six specific tags:

- <mdoc></mdoc>, to open a documentation field in the FAUST code,
 - <equation></equation>, to get equations of a FAUST expression,
 - <diagram></diagram>, to get the top-level block-diagram of a FAUST expression,
 - <metadata></metadata>, to reference FAUST metadatas (cf. declarations), calling the corresponding keyword,
 - <notice />, to insert the "adaptive" notice all formulas actually printed,
 - - (listing [attributes] />, to insert the listing of FAUST files called.

The tag can have up to three boolean attributes (set to "true" by default):

- mdoctags for <mdoc> tags;
- · dependencies for other files dependencies;
- distributed for the distribution of interleaved FAUST code between <mdoc> sections.

12.5.2 The mdoc top-level tags

The <mdoc></mdoc> tags are the top-level delimiters for FAUST mathematical documentation sections. This means that the four other documentation tags can't be used outside these pairs (see section 3.2.3).

In addition of the four inner tags, <mdoc></mdoc> tags accept free LaTEX text, including its standard macros (like \section, \emph, etc.). This allows to manage the presentation of resulting tex file directly from within the input FAUST file.

The complete list of the LATEX packages included by FAUST can be found in the file architecture/latexheader.tex.

12.5.3 An example of manual mathdoc

```
\title{<metadata>name</metadata>}
\author{<metadata>author</metadata>}
\date{\today}
\maketitle
\begin{tabular}{11}
   \hline
   \textbf{copyright} & <metadata>copyright</metadata>
   \hline
\end{tabular}
\bigskip
</mdoc>
//
// Noise generator and demo file for the Faust math
  documentation
//
declare name "Noise";
declare version
                 "1.1";
                 "Grame";
declare author
declare author
                 "Yghe";
                "BSD";
declare license
declare copyright "(c)GRAME 2009";
<mdoc>
\section{Presentation of the "noise.dsp" Faust program}
This program describes a white noise generator with an
interactive volume, using a random function.
\subsection{The random function}
</mdoc>
random = +(12345)^**(1103515245);
```

```
<mdoc>
The \texttt{random} function describes a generator of
  random numbers, which equation follows. You should
  notice hereby the use of an integer arithmetic on 32
  bits, relying on integer wrapping for big numbers.
<equation>random</equation>
\subsection{The noise function}
</mdoc>
noise = random/2147483647.0;
The white noise then corresponds to:
<equation>noise</equation>
\subsection{Just add a user interface element to play
  volume!}
</mdoc>
process = noise * vslider("Volume[style:knob]", 0, 0, 1,
  0.1);
< mdoc >
Endly, the sound level of this program is controlled by a
   user slider, which gives the following equation:
<equation>process</equation>
\section{Block-diagram schema of process}
This process is illustrated on figure 1.
<diagram>process</diagram>
\section{Notice of this documentation}
You might be careful of certain information and naming
  conventions used in this documentation:
<notice />
\section{Listing of the input code}
The following listing shows the input Faust code, parsed
  to compile this mathematical documentation.
<listing mdoctags="false" dependencies="false"</pre>
distributed="true" />
```

</mdoc>

The following page which gathers the four resulting pages of noise.pdf in small size. might give you an idea of the produced documentation.

12.5.4 The -stripmdoc option

As you can see on the resulting file noisemetadata.pdf on its pages 3 and 4, the listing of the input code (section 4) contains all the mathdoc text (here colored in grey). As it may be useless in certain cases (see Goals, section 12.1), we provide an option to strip mathdoc contents directly at compilation stage: -stripmdoc (short) or --strip-mdoc -tags (long).

12.6 Localization of mathdoc files

By default, texts used by the documentator are in English, but you can specify another language (French, German and Italian for the moment), using the -mdlang (or --mathdoc-lang) option with a two-letters argument (en, fr, it, etc.).

The faust2mathdoc script also supports this option, plus a third short form with -1:

```
faust2mathdoc -l fr myfaustfile.dsp
```

If you would like to contribute to the localization effort, feel free to translate the mathdoc texts from any of the mathdoctexts-*.txt files, that are in the architecture directory (mathdoctexts-fr.txt, mathdoctexts-it.txt, etc.). As these files are dynamically loaded, just adding a new file with an appropriate name should work.

Noise

Grame, Yghe

March 9, 2010

name
version
author
license
copyright Noise
1.1
Grame, Yghe
BSD
BSD
(c)GRAME 2009

declare name "Noise";
declare version "1.1";
declare author "Grame";
declare author "Rghe';
declare 1conne "ESD';
declare 1conne "ESD'; // Noise generator and demo file for the Faust math documentation

1 Presentation of the "noise.dsp" Faust program

This program describes a white noise generator with an interactive volume, using a random function.

1.1 The random function

random = +(int(12345))~*(int(1103515245));

The random function describes a generator of random numbers, which equation follows. You should notice hereby the use of an integer arithmetic on 32 bits, relying on integer varapping for big numbers.

1. Output signal y such that

 $y(t) = r_1(t)$

Input signal (none)

3. Intermediate signal r_1 such that

 $r_1(t) = 12345 \oplus 1103515245 \odot r_1(t\!-\!1)$

1.2 The noise function

The white noise then corresponds to: 1. Output signal y such that

noise = (int(random))/(int(random+1));

 $y(t) = s_1(t)$

Input signal (none)

3. Intermediate signal s_1 such that

 $s_1(t) = \operatorname{int}\left(r_1(t)\right) \oslash \operatorname{int}\left(1 \oplus r_1(t)\right)$

1.3 Just add a user interface element to play volume!

Endly, the sound level of this program is controlled by a user slider, which gives the following equation:

process = noise * vslider("Volume[style:knob]", 0, 0, 1, 0.1);

 $y(t) = u_{s1}(t) \cdot s_1(t)$

Input signal (none)

3. User-interface input signal u_{s1} such that

"Volume" $u_{s1}(t) \in [0,1]$ (default value = 0)

2 Block-diagram schema of process

This process is illustrated on figure 1.

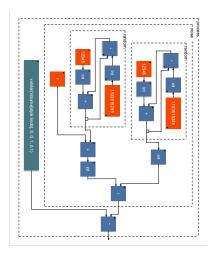


Figure 1: Block diagram of process

3 Notice of this documentation

You might be careful of certain information and naming conventions used in

- \bullet This document was generated using Faust version 0.9.13 on March 09, 2010.
- The value of a Faust program is the result of applying the signal trans-former denoted by the expression to which the process identifier is bound to input signals, running at the fs sampling frequency.
- Faust (Functional Audio Stream) is a functional programming language designed for synchronous real-time signal processing and synthesis appli-cations. A Faust program is a set of bindings of identifiers to expression that denote signal transformers. A signal s in S is a function mapping! times t∈ Z to values s(t) ∈ R, while a signal transformer is a function

from S^n to $S^m,$ where $n,m\in\mathbb{N}.$ See the Faust manual for additional information (http://faust.grame.fr).

- Every mathematical formula derived from a Faust expression is assumed. in this document, to having been normalized (in an implementation-dependent manner) by the Faust compiler.
- A block diagram is a graphical representation of the Faast binding of an identifier I to an expression E; each graph is put in a box labeled by I. Subexpressions of E are recursively displayed as long as the whole picture fits in one page.
- ∀x ∈ ℝ,

$$\operatorname{int}(x) = \left\{ \begin{array}{ll} \lfloor x \rfloor & \text{if } x > 0 \\ \lceil x \rceil & \text{if } x < 0 \\ 0 & \text{if } x = 0 \end{array} \right..$$

 $\bullet\,$ This document uses the following integer operations:

operation	name	semantics
$i \oplus j$	integer addition	$normalize(i + j)$, in \mathbb{Z}
a. ⊙ j.	integer multiplication	$normalize(i \cdot j), \text{ in } \mathbb{Z}$
$i \oslash j$	integer division	normalize(int (i/j)), in \mathbb{Q}

Integer operations in Faust are inspired by the semantics of operations on the n-bit two's complement representation of integer numbers; they are internal composition laws on the subset $[-2^{n-1}, 2^{n-1}-1]$ of \mathbb{Z} , with n=32. For any integer binary operation \times on \mathbb{Z} , the \otimes operation is defined as: $i\otimes j=$ normalize($i\times j$), with

$$\operatorname{normalize}(i) = i - N \cdot \operatorname{sign}(i) \cdot \left\lfloor \frac{|i| + N/2 + (\operatorname{sign}(i) - 1)/2}{N} \right\rfloor$$

operations are defined likewise. where $N=2^n$ and sign(i)=0 if i=0 and i/|i| otherwise. Unary integer

- The noisemetadata-mdoc/ directory may also include the following sub-
- cpp/ for Faust compiled code;
- pdf/ which contains this document;
- src/ for all Faust sources used (even libraries);
- svg/ for block diagrams, encoded using the Scalable Vector Graphics format (http://www.w3.org/Graphics/SVG/);
- tex/ for the EMEX source of this document.

¹Faust assumes that $\forall s \in S, \forall t \in \mathbb{Z}, s(t) = 0$ when t < 0.

4 Listing of the input code

The following listing shows the input Faust code, parsed to compile this mathematical documentation.

Listing 1: noisemetadata.dsp

Ü

12.7 Summary of the mathdoc generation steps

- I. First, to get the full mathematical documentation done on your faust file, call faust2mathdoc myfaustfile.dsp.
- 2. Then, open the pdf file myfaustfile-mdoc/pdf/myfaustfile.pdf.
- 3. That's all!

Chapter 13

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